# Products of Factorials in Smarandache Type Expressions

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## Introduction

In [3] and [5] the authors ask how many primes are of the form  $x^y + y^x$ , where gcd (x, y) = 1 and  $x, y \ge 2$ . Moreover, Jose Castillo (see [2]) asks how many primes are of the Smarandache form  $x_1^{x_2} + x_2^{x_3} + ... + x_n^{x_1}$ , where n > 1,  $x_1$ ,  $x_2$ , ...,  $x_n > 1$  and gcd  $(x_1, x_2, ..., x_n) = 1$  (see [9]).

In this article we announce a lower bound for the size of the largest prime divisor of an expression of the type  $ax^y + by^x$ , where  $ab \neq 0$ ,  $x, y \geq 2$  and gcd (x, y) = 1.

For any finite extension F of Q let  $d_F = [F : Q]$ . For any algebraic number  $\zeta \in F$  let  $N_F(\zeta)$  denote the norm of  $\zeta$ .

For any rational integer n let P(n) be the largest prime number P dividing n with the convention that  $P(0) = P(\pm 1) = 1$ .

Theorem 1. Let  $\alpha$  and  $\beta$  be algebraic integers with  $\alpha \cdot \beta \neq 0$ . Let  $K = \mathbb{Q}[\alpha, \beta]$ . For any two positive integers x and y let  $X = \max(x, y)$ . There exist computable positive numbers  $C_1$  and  $C_2$  depending only on  $\alpha$  and  $\beta$  such that

$$P\left(N_{\mathbf{K}}\left(\alpha x^{y} + \beta y^{x}\right)\right) > C_{1}\left(\frac{X}{\log^{3} X}\right)^{1/(d_{\mathbf{K}}+1)}$$

whenever  $x, y \ge 2$ , gcd(x, y) = 1, and  $X > C_2$ .

The proof of Theorem 1 uses lower bounds for linear forms in logarithms of algebraic numbers (see [1] and [7]) as well as an idea of Stewart (see [10]).

Erdös and Obláth (see [4]) found all the solutions of the equation  $n! = x^p \pm y^p$  with gcd (x, y) = 1 and p > 2. Moreover, the author (see [6]) showed that in every non-degenerate binary recurrence sequence  $(u_n)_{n\geq 0}$  there are only finitely many terms which are products of factorials.

We use Theorem 1 to show that for any two given integers a and b with  $ab \neq 0$ , there exist only finitely many numbers of the type  $ax^y + by^x$ , where  $x, y \geq 2$  and gcd (x, y) = 1, which are products of factorials.

Let PF be the set of all positive integers which can be written as products of factorials; that is

$$\mathcal{PF} = \{ w \mid w = \prod_{j=1}^k m_j!, \text{ for some } m_j \geq 1 \}.$$

Theorem 2. Let  $f_1, ..., f_s \in \mathbf{Z}[X, Y]$  be  $s \geq 1$  homogeneous polynomials of positive degrees. Assume that  $f_i(0, Y) \cdot f_i(X, 0) \not\equiv 0$  for i = 1, ..., s. Then, the equation

$$f_1(x_1^{y_1}, y_1^{x_1}) \cdot \dots \cdot f_s(x_s^{y_s}, y_s^{x_s}) \in \mathcal{PF},$$
 (1)

with gcd  $(x_i, y_i) = 1$  and  $x_i, y_i \geq 2$ , for i = 1, ..., s, has finitely many solutions  $x_1, y_1, ..., x_s, y_s$ . Moreover, there exists a computable positive number C depending only on the polynomials  $f_1, ..., f_s$  such that all solutions of equation (1) satisfy  $\max(x_1, y_1, ..., x_s, y_s) < C$ .

We also have the following inhomogeneous variant of theorem 2.

Theorem 3. Let  $f_1, ..., f_s \in \mathbb{Z}[X]$  be  $s \geq 1$  polynomials of positive degrees. Assume that  $f_i(0) \equiv 1 \pmod 2$  for i = 1, ..., s. Let  $a_1, ..., a_s$  and  $b_1, ..., b_s$  be 2s odd integers. Then, the equation

$$f_1(a_1x_1^{y_1} + b_1y_1^{x_1}) \cdot \dots \cdot f_s(a_sx_s^{y_s} + b_sy_s^{x_s}) \in \mathcal{PF},$$
 (2)

with gcd  $(x_i, y_i) = 1$  and  $x_i, y_i \geq 2$ , for i = 1, ..., s, has finitely many solutions  $x_1, y_1, ..., x_s, y_s$ . Moreover, there exists a computable positive number C depending only on the polynomials  $f_1, ..., f_s$  and the 2s numbers  $a_1, b_1, ..., a_s, b_s$ , such that all solutions of equation (2) satisfy  $\max(x_1, y_1, ..., x_s, y_s) < C$ .

We conclude with the following computational results:

Theorem 4. All solutions of the equation

$$x^y \pm y^x \in \mathcal{PF}$$
 with gcd  $(x, y) = 1$  and  $x, y \ge 2$ ,

satisfy  $\max(x, y) < \exp 177$ .

Theorem 5. All solutions of the equation

$$x^{y} + y^{z} + z^{x} = n!$$
 with gcd  $(x, y, z) = 1$  and  $x, y, z \ge 2$ ,

satisfy max  $(x, y, z) < \exp 518$ .

### 2. Preliminary Results

The proofs of theorems 1-5 use estimations of linear forms in logarithms of algebraic numbers.

Suppose that  $\zeta_1$ , ...,  $\zeta_l$  are algebraic numbers, not 0 or 1, of heights not exceeding  $A_1$ , ...,  $A_l$ , respectively. We assume  $A_m \geq e^e$  for m = 1, ..., l. Put  $\Omega = \log A_1 ... \log A_l$ . Let  $F = \mathbb{Q}[\zeta_1, ..., \zeta_l]$ . Let  $n_1, ..., n_l$  be integers, not all 0, and let  $B \geq \max |n_m|$ . We assume  $B \geq e^2$ . The following result is due to Baker and Wüstholz.

Theorem BW ([1]). If  $\zeta_1^{n_1}...\zeta_m^{n_l} \neq 1$ , then

$$|\zeta_1^{n_1}...\zeta_l^{n_l} - 1| > \frac{1}{2} \exp(-(16(l+1)d_{\mathbf{F}})^{2(l+3)}\Omega \log B).$$
 (3)

In fact, Baker and Würtholz showed that if  $\log \zeta_1$ , ...,  $\log \zeta_l$  are any fixed values of the logarithms, and  $\Lambda = n_1 \log \zeta_1 + ... + n_l \log \zeta_l \neq 0$ , then

$$\log |\Lambda| > -(16ld_{\mathbf{F}})^{2(l+2)} \Omega \log B. \tag{4}$$

Now (4) follows easily from (3) via an argument similar to the one used by Shorey et al. in their paper [8].

We also need the following p-adic analogue of theorem BW which is due to van der Poorten.

Theorem vdP ([7]). Let  $\pi$  be a prime ideal of F lying above a prime integer p. Then,

$$\operatorname{ord}_{\pi}(\zeta_{1}^{n_{1}}...\zeta_{l}^{n_{l}}-1) < (16(l+1)d_{\mathbf{F}})^{12(l+1)} \frac{p^{d_{\mathbf{F}}}}{\log p} \Omega(\log B)^{2}.$$
 (5)

The following estimations are useful in what follows.

Lemma 1. Let  $n \ge 2$  be an integer, and let  $p \le n$  be a prime number. Then

$$n^{n/2} \le n! \le n^n. \tag{6}$$

(ii) 
$$\frac{n}{4(p-1)} \le \operatorname{ord}_p n! \le \frac{n}{p-1}.$$
 (7)

Proof. See [6].

Lemma 2. (1) Let  $s \ge 1$  be a positive integer. Let C and X be two positive numbers such that  $C > \exp s$  and X > 1. Let y > 0 be such that  $y < C \log^s X$ . Then,  $y \log y < (C \log C) \log^{s+1} X$ .

(2) Let  $s \ge 1$  be a positive integer, and let  $C > \exp(s(s+1))$ . If X is a positive number such that  $X < C \log^s X$ , then  $X < C \log^{s+1} C$ .

Proof. (1) Clearly,

$$y \log y < C \log^s X (\log C + s \log \log X).$$

It suffices to show that

$$\log C + s \log \log X < \log C \log X$$
.

The above inequality is equivalent to

$$\log C(\log X - 1) > s \log \log X.$$

This last inequality is obviously satisfied since  $\log C > s$  and  $\log X > \log \log X + 1$ , for all X > 1.

(2) Suppose that  $X \ge C \log^{s+1} C$ . Since  $s \ge 1$  and  $C > \exp(s(s+1))$ , it follows that  $C \log^{s+1} C > C > \exp s$ . The function  $\frac{y}{\log^s y}$  is increasing for  $y > \exp s$ . Hence, since  $X \ge C \log^{s+1} C$ , we conclude that

$$\frac{C\log^{s+1}C}{\log^s(C\log^{s+1}C)} \le \frac{X}{\log^sX} < C.$$

The above inequality is equivalent to

$$\frac{\log^{s+1} C}{\left(\log C + (s+1)\log\log C\right)^s} < 1,$$

or

$$\log C < \left(1 + (s+1)\frac{\log\log C}{\log C}\right)^{s}.$$

By taking logarithms in this last inequality we obtain

$$\log \log C < s \log \left(1 + (s+1) \frac{\log \log C}{\log C}\right) < s(s+1) \frac{\log \log C}{\log C}.$$

This last inequality is equivalent to  $\log C < s(s+1)$ , which contradicts the fact that  $C > \exp(s(s+1))$ .

### 3. The Proofs

The Proof of Theorem 1. By  $C_1$ ,  $C_2$ , ..., we shall denote computable positive numbers depending only on the numbers  $\alpha$  and  $\beta$ . Let  $d = d_K$ . Let

$$N_{\mathbf{K}}(\alpha x^y + \beta y^x) = p_1^{\delta_1} \cdot \dots \cdot p_k^{\delta_k}$$

where  $2 < p_1 < p_2 < ... < p_k$  are prime numbers. For  $\mu = 1, ..., d$ , let  $\alpha^{(\mu)}x^y + \beta^{(\mu)}y^x$  be a conjugate, in K, of  $\alpha x^y + \beta y^x$ . Fix i = 1, ..., k. Let  $\pi$  be a prime ideal of K lying above  $p_i$ . We use theorem vdP to bound  $\operatorname{ord}_{\pi}(\alpha^{(\mu)}x^y + \beta^{(\mu)}y^x)$ . We distinguish two cases:

CASE 1.  $p_i \mid xy$ . Suppose, for example, that  $p_i \mid y$ . Since (x, y) = 1, it follows that  $p_i \not\mid x$ . Hence, by theorem vdP,

$$\operatorname{ord}_{\pi}\left(\alpha^{(\mu)}x^{y} + \beta^{(\mu)}y^{x}\right) = \operatorname{ord}_{\pi}\left(\alpha^{(\mu)}x^{y}\right) + \operatorname{ord}_{\pi}\left(1 - \left(-\frac{\beta^{(\mu)}}{\alpha^{(\mu)}}\right)y^{x}x^{-y}\right) < 0$$

$$< C_1 + C_2 \frac{p_i^d}{\log p_i} \log^4 X.$$
 (8)

where  $C_1 = d \cdot \log_2 N_{\mathbf{K}}(\alpha)$ , and  $C_2$  can be computed in terms of  $\alpha$  and  $\beta$  using theorem vdP.

CASE 2.  $p_i \nmid xy$ . In this case

$$\operatorname{ord}_{\pi}\left(\alpha^{(\mu)}x^{y} + \beta^{(\mu)}y^{x}\right) = \operatorname{ord}_{\pi}\left(\alpha^{(\mu)}x^{y}\right) + \operatorname{ord}_{\pi}\left(1 - \left(-\frac{\beta^{(\mu)}}{\alpha^{(\mu)}}\right) \cdot \frac{y^{x}}{x^{y}}\right) < 0$$

$$< C_1 + C_2 \frac{p_i^d}{\log p_i} \log^4 X.$$
 (9)

Combining Case 1 and Case 2 we conclude that

$$\operatorname{ord}_{\pi} \left( \alpha^{(\mu)} x^{y} + \beta^{(\mu)} y^{x} \right) < C_{3} \frac{p_{i}^{d}}{\log p_{i}} \log^{4} X, \tag{10}$$

where  $C_3 = 2 \cdot \max (C_1, C_2)$ . Hence,

$$\delta_i = \operatorname{ord}_{p_i} \left( N_{\mathbf{K}} \left( \alpha x^y + \beta y^x \right) \right) < C_4 \frac{p_i^d}{\log p_i} \log^4 X.$$
 (11)

where  $C_4 = dC_3$ . Denote  $p_k$  by P. Since  $p_i \leq P$  for i = 1, ..., k, it follows, by formula (11), that

$$\log \left( N_{\mathbf{K}} \left( \alpha x^{y} + \beta y^{x} \right) \right) \leq \sum_{i=1}^{k} \delta_{i} \cdot \log p_{i} < kC_{4} P^{d} \log^{4} X.$$
 (12)

Clearly  $k \leq \pi(P)$ , where  $\pi(P)$  is the number of primes less than or equal to P. Combining inequality (12) with the prime number theorem we conclude that

$$\log\left(N_{\mathbf{K}}(\alpha x^{y} + \beta y^{x})\right) < C_{5} \frac{P^{d+1}}{\log P} \log^{4} X. \tag{13}$$

We now use theorem BW to find a lower bound for  $\log(N_K(\alpha x^y + \beta y^x))$ . Suppose that X = y. For  $\mu = 1, ..., d$ , we have

$$\log\left(\left|\alpha^{(\mu)}x^{y} + \beta^{(\mu)}y^{x}\right|\right) = \log\left(\left|\alpha^{(\mu)}x^{y}\right|\right) + \log\left(\left|1 - \left(-\frac{\beta^{(\mu)}}{\alpha^{(\mu)}}\right) \frac{y^{x}}{x^{y}}\right|\right) >$$

$$> C_{6} + X\log 2 - C_{7}\log^{3}X.$$

where  $C_6 = \min \left( \log |\alpha^{(\mu)}| \mid \mu = 1, ..., d \right)$ , and  $C_7$  can be computed using theorem BW. Hence,

$$\log\left(N_{\mathbf{K}}(\alpha x^{y} + \beta y^{x})\right) > dC_{6} + dX \log 2 - dC_{7} \log^{3} X. \tag{14}$$

Let  $C_8 = dC_6$ ,  $C_9 = d \log 2$ , and  $C_{10} = dC_7$ . Let also  $C_{11}$  be the smallest positive number such that

$$\frac{1}{2}C_9 y > C_{10} \log^3 y - C_8, \qquad \text{for } y > C_{11}.$$

Combining inequalities (13) and (14) it follows that

$$C_5 \frac{P^{d+1}}{\log P} \log^4 X > C_8 + C_9 X - C_{10} \log^3 X > \frac{1}{2} C_9 X, \tag{15}$$

for  $X \ge C_{11}$ . Inequality (15) clearly shows that

$$P > C_{12} \left( \frac{X}{\log^3 X} \right)^{\frac{1}{d+1}}, \quad \text{for } X \ge C_{11}.$$

The Proof of Theorem 2. By  $C_1$ ,  $C_2$ , ..., we shall denote computable positive numbers depending only on the polynomials  $f_1$ , ...,  $f_s$ . We may assume that  $f_1$ , ...,  $f_s$  are linear forms with algebraic coefficients. Let  $f_i(X, Y) = \alpha_i X + \beta_i Y$  where  $\alpha_i \beta_i \neq 0$ , and let  $K = \mathbb{Q}[\alpha_1, \beta_1, ..., \alpha_s, \beta_s]$ . Let  $(x_1, y_1, ..., x_s, y_s)$  be a solution of (1). Equation (1) implies that

$$\prod_{i=1}^{s} N_{\mathbf{K}} \left( \alpha_i x_i^{y_i} + \beta_i y_i^{x_i} \right) = n_1! \cdot \dots \cdot n_k! \tag{16}$$

We may assume that  $2 \le n_1 \le n_2 \le ... \le n_k$ . Let  $X = \max(x_i, y_i \mid i = 1, ..., s)$ . It follows easily, by inequality (10), that

$$\operatorname{ord}_{2}\left(\prod_{i=1}^{s} N_{K}\left(\alpha_{i} x_{i}^{y_{i}} + \beta_{i} y_{i}^{x_{i}}\right)\right) < C_{1} \log^{4} X.$$

$$(17)$$

Hence,

$$\sum_{i=1}^k \operatorname{ord}_2 n_i! < C_1 \log^4 X.$$

By lemma 1, it follows that

$$n_k < 4C_1 \log^4 X. \tag{18}$$

On the other hand, by theorem 1, there exists computable constants  $C_{2i}$  and  $C_{3i}$ , such that

$$P\left(N_{\mathbf{K}}\left(\alpha_{i}x_{i}^{y_{i}}+\beta_{i}y_{i}^{x_{i}}\right)\right) > C_{2i}\left(\frac{X_{i}}{\log^{3}X_{i}}\right)^{1/(d_{\mathbf{K}}+1)}$$

$$\tag{19}$$

whenever  $x_i$ ,  $y_i \ge 2$ ,  $gcd(x_i, y_i) = 1$  and  $X_i = max(x_i, y_i) > C_{3i}$ . Let  $C_2 = min(C_{2i} \mid i = 1, ..., s)$  and let  $C_3 = max(C_{3i} \mid i = 1, ..., s)$ . Suppose that  $X > C_3$ . From inequality (19) we conclude that

$$P\left(\prod_{i=1}^{s} N_{\mathbf{K}}\left(\alpha_{i} x_{i}^{y_{i}} + \beta_{i} y_{i}^{x_{i}}\right)\right) > C_{2}\left(\frac{X}{\log^{3} X}\right)^{1/(d_{\mathbf{K}}+1)}.$$
 (20)

Since  $P \mid \prod_{i=1}^{k} n_i!$ , it follows that  $P \leq n_k$ . Combining inequalities (18) and (20) we conclude that

$$C_2 \left(\frac{X}{\log^3 X}\right)^{1/(d_{\mathbf{K}}+1)} < 4C_1 \log^4 X.$$
 (21)

Inequality (21) clearly shows that  $X < C_4$ .

The Proof of Theorem 3. By  $C_1$ ,  $C_2$ , ..., we shall denote computable positive numbers depending only on the polynomials  $f_1$ , ...,  $f_s$  and on the numbers  $a_1$ ,  $b_1$ , ...,  $a_s$ ,  $b_s$ . Let  $(x_1, y_1, ..., x_s, y_s)$  be a solution of (2). Let  $X_i = \max(x_i, y_i)$ , and let  $X = \max(X_i \mid i = 1, ..., s)$ . Finally, let

$$f_i(Z) = c_i \prod_{j=1}^{d_i} (Z - \zeta_{i,j}).$$

Let  $K = \mathbb{Q}[\zeta_{i,j}]_{\substack{1 \le i \le s \\ 1 \le j \le d_i}}^{1 \le i \le s}$ , and let  $d = [K : \mathbb{Q}]$ ,  $D = \sum_{i=1}^{s} d_i$ , and  $c = \prod_{i=1}^{s} c_i$ .

Let  $\pi$  be a prime ideal of K lying above 2. Let  $Z_i = a_i x_i^{y_i} + b_i y_i^{x_i}$ . We first bound  $\operatorname{ord}_{\pi} f_i(Z_i)$ . First, notice that  $\operatorname{ord}_{\pi} (a_i b_i) = 0$ . Moreover, since  $f_i(0) \equiv 1 \pmod{2}$ , it follows that  $\operatorname{ord}_{\pi} (\zeta_{i,j}) = 0$ , for all  $j = 1, ..., d_i$ . We distinguish 2 cases:

CASE 1. Assume that  $2 \not\mid x_i y_i$ . Then  $f_i(Z_i) \equiv f_i(0) \equiv 1 \pmod{2}$ . Hence,  $\operatorname{ord}_{\pi} f_i(Z_i) = 0$ .

CASE 2. Assume that  $2 \mid x_i$ . In this case,  $\operatorname{ord}_{\pi}(y) = 0$ . Fix  $j = 1, ..., d_i$ . Then,

$$\operatorname{ord}_{\pi}\left(Z_{i}-\zeta_{i,j}\right)=\operatorname{ord}_{\pi}\left(a_{i}x_{i}^{y_{i}}+\left(b_{i}y_{i}^{x_{i}}-\zeta_{i,j}\right)\right). \tag{22}$$

Since  $\operatorname{ord}_{\pi}(b_i y_i^{x_i}) = \operatorname{ord}_{\pi}(\zeta_{i,j}) = 0$ , it follows, by theorem vdP, that

$$\operatorname{ord}_{\pi}(b_{i}y_{i}^{x_{i}} - \zeta_{i,j}) = \operatorname{ord}_{\pi}(b_{i}y_{i}^{x_{i}}(\zeta_{i,j})^{-1} - 1) < C_{1}\log^{3}X_{i}.$$
 (23)

We distinguish 2 cases:

CASE 2.1.  $y_i \ge C_1 \log^3 X_i$ . In this case, from formula (22) and inequality (23), it follows that

$$\operatorname{ord}_{\pi}(Z_{i} - \zeta_{i,j}) = \operatorname{ord}_{\pi}(b_{i}y_{i}^{x_{i}} - \zeta_{i,j}) < C_{1} \log^{3} X_{i}. \tag{24}$$

CASE 2.2.  $y_i < C_1 \log^3 X_i$ . In this case,

$$\operatorname{ord}_{\pi}(Z_{i} - \zeta_{i,j}) = \operatorname{ord}_{\pi}\left(b_{i}y_{i}^{x_{i}} + \left(a_{i}x_{i}^{y_{i}} - \zeta_{i,j}\right)\right). \tag{25}$$

Let  $\Delta = a_i x_i^{y_i} - \zeta_{i,j}$ . Let  $H(\Delta)$  be the height of  $\Delta$ . Clearly,

$$H(\Delta) < C_2 x_i^{d_i y_i}.$$

Hence,

$$\log(H(\Delta)) < \log C_2 + d_i y_i \log x_i < C_3 + C_4 \log^4 X_i,$$

where  $C_3 = \log C_2$ , and  $C_4 = C_1 \cdot \max(d_i \mid i = 1, ..., s)$ . Since  $\operatorname{ord}_{\pi}(b_i) = \operatorname{ord}_{\pi}(y_i^{x_i}) = 0$ , it follows, by theorem vdP, that

$$\operatorname{ord}_{\pi}(Z_{i} - \zeta_{i,j}) = \operatorname{ord}_{\pi}(1 - b_{i}^{-1}y_{i}^{-x_{i}}\Delta) < C_{5}\log y_{i}\log(H(\Delta))\log^{2}x_{i} < C_{5}\log y_{i}\log(H(\Delta))$$

$$< C_5 \log^3 X_i (C_3 + C_4 \log^4 X_i).$$
 (26)

Let  $C_6 = 2C_4C_5$ . Also, let

$$C_7 = \exp((C_3/C_4)^{1/4}).$$

From inequalities (23) and (26), it follows that

$$\operatorname{ord}_{\pi}(Z_{i} - \zeta_{i,j})) < C_{6} \log^{7} X, \qquad \text{for } X > C_{7}. \tag{27}$$

Hence,

$$\operatorname{ord}_{2}\left(\prod_{i=1}^{s} f_{i}(Z_{i})\right) < C_{8} \log^{7} X, \qquad \text{for } X > C_{7}, \tag{28}$$

where  $C_8 = 2\max (sDC_6, c)$ . Suppose now that

$$\prod_{i=1}^{s} f_i(Z_i) = \prod_{j=1}^{k} n_j!, \tag{29}$$

where  $2 \le n_1 \le n_2 \le ... \le n_k$ . From inequality (28) and lemma 1, it follows that

$$\sum_{j=1}^k n_j < C_9 \log^7 X,$$

where  $C_9 = 4C_8$ . Hence,

$$\log\left(\prod_{j=1}^{k} n_{j}!\right) = \sum_{j=1}^{k} \log n_{j}! < \sum_{j=1}^{k} n_{j} \log n_{j} < \left(\sum_{j=1}^{k} n_{j}\right) \log\left(\sum_{j=1}^{k} n_{j}\right) < < C_{9} \log^{7} X \left(\log C_{9} + 7 \log \log X\right), \quad \text{for } X > C_{7}.$$
 (30)

Let  $C_{10}$  be the smallest positive number  $\geq C_7$  such that

$$y > \log C_9 + 7 \log \log y, \qquad \text{for } y > C_{10}.$$

From inequality (30), it follows that

$$\log\left(\prod_{j=1}^{k} n_{j}!\right) < C_{9} \log^{8} X, \qquad \text{whenever } X > C_{10}. \tag{31}$$

We now bound  $\log(\prod_{i=1}^{s} f_i(Z_i))$ . Fix i = 1, ..., s. Suppose that  $y_i = X_i$ . By Theorem BW,

$$\log |Z_{i}| = \log |a_{i}x_{i}^{y_{i}} + b_{i}y_{i}^{x_{i}}| = \log(|a_{i}|x_{i}^{y_{i}}) + \log\left(\left|1 - \left(-\frac{b_{i}}{a_{i}}\right)y_{i}^{x_{i}}x_{i}^{-y_{i}}\right|\right) >$$

$$> C_{11} + X_{i} \log 2 - C_{12} \log^{3} X_{i},$$
(32)

where  $C_{11} = \min (|a_i| | i = 1, ..., s)$ , and  $C_{12}$  can be computed using theorem BW. Let  $C_{13} = (\log 2)/2$ , and let  $C_{14}$  be the smallest positive number  $\geq C_{10}$  such that

$$C_{11} + y \log 2 - C_{12} \log^3 y > C_{13}y,$$
 for  $y > C_{14}$ .

From inequality (32) it follows that

$$\max (\log |Z_i|) > C_{13}X, \quad \text{for } X > C_{14}.$$
 (33)

On the other hand, for each i = 1, ..., s, there exists two computable constants  $C_i$  and  $C'_i$  such that

$$|f_i(Z_i)| > C_i |Z_i|^{d_i}$$
, whenever  $|Z_i| > C_i'$ .

Let  $C_{15} = \min (C_i \mid i = 1, ..., s)$ , and let  $C_{16} = \max (C_i' \mid i = 1, ..., s)$ . Finally, let  $C_{17} = \max (C_{14}, (\log C_{16})/C_{13})$ . Suppose that  $X > C_{17}$ . Since  $|f_i(Z_i)| \ge 1$ , for all i = 1, ..., s, it follows, by inequality (33), that

$$\log \left(\prod_{i=1}^s f_i(Z_i)\right) \geq \max \left(\log |f_i(Z_i)| \ i=1, \ ..., \ s\right) >$$

$$> \log C_{15} + \max \left( \log |Z_i| \mid i = 1, ..., s \right) > \log C_{15} + C_{13}X, \text{ for } X > C_{17}.$$
(34)

From equation (29) and inequalities (31) and (34), it follows that

$$\log C_{15} + C_{13}X < C_9 \log^8 X,$$
 for  $X > C_{17}$ . (35)

Inequality (35) clearly shows that  $X < C_{18}$ .

The Proof of Theorem 4. Let  $X = \max(x, y)$ . Notice that if  $x^y \pm y^x \in \mathcal{PF}$ , than xy is odd. Hence, by theorem vdP,

$$\operatorname{ord}_{2}(x^{y} \pm y^{x}) = \operatorname{ord}_{2}(1 - (\mp y)^{x}x^{-y}) < 48^{36} \cdot \frac{2}{\log 2} \cdot \log^{4} X.$$
 (36)

Suppose that

$$x^y \pm y^x = n_1! \cdot \dots \cdot n_k!, \tag{37}$$

where  $2 \le n_1 \le ... \le n_k$ . From inequality (36) and lemma 1 it follows that

$$\sum_{i=1}^{k} n_i \le 4 \left( \sum_{i=1}^{k} \operatorname{ord}_2(n_i!) \right) < 48^{36} \cdot \frac{8}{\log 2} \cdot \log^4 X < 12 \cdot 48^{36} \cdot \log^4 X. \tag{38}$$

It follows, by lemma 2 (1), that

$$\log(x^y \pm y^x) = \log \prod_{i=1}^k n_i! = \sum_{i=1}^k \log n_i! < \sum_{i=1}^k n_i \log n_i <$$

$$< \left(\sum_{i=1}^{k} n_i\right) \log \left(\sum_{i=1}^{k} n_i\right) < 12 \cdot 48^{36} \log \left(12 \cdot 48^{36}\right) \cdot \log^5 X < 1703 \cdot 48^{36} \log^5 X. \tag{39}$$

Suppose now that X = y. Then, by theorem BW,

$$\log|x^{y} \pm y^{x}| \ge \log|x^{y} - y^{x}| = \log(x^{y}) + \log|1 - y^{x}x^{-y}| >$$

$$> X \log 3 - \log 2 - 48^{10} \log^{3} X.$$
(40)

Combining inequalities (39) and (40), we conclude that

$$X < X \log 3 < \log 2 + 48^{10} \log^3 X + 1703 \cdot 48^{36} \log^5 X < 1704 \cdot 48^{36} \log^5 X. \tag{41}$$

Let  $C = 1704 \cdot 48^{36}$ , and let s = 5. Since  $\log C = \log 1704 + 36 \log 48 > 30$ , it follows, by lemma 2 (2), that

$$X < C \cdot \log^6 C < 1704 \cdot 48^{36} \cdot 147^6. \tag{42}$$

Hence,  $\log X < 177$ .

The Proof of Theorem 5. Suppose that (x, y, z, n) is a solution of  $x^y + y^z + z^x = n!$ , with gcd (x, y, z) = 1 and min (x, y, z) > 1. Let  $X = \max(x, y, z)$ . We assume that  $\log X > 519$ . Clearly, not all three numbers x, y, z can be odd. We may assume that  $2 \mid x$ . In this case, both y and z are odd. By theorem vdP,

$$\operatorname{ord}_{2}(y^{z} + z^{x}) = \operatorname{ord}_{2}(1 - (-y)^{-z}z^{x}) < 48^{36} \frac{2}{\log 2} \log^{4} X < 3 \cdot 48^{36} \log^{4} X.$$
(43)

We distinguish two cases:

CASE 1.  $y \ge 3 \cdot 48^{36} \log^4 X$ . In this case, by lemma 1,

$$n/4 \le \operatorname{ord}_2 n! = \operatorname{ord}_2(x^y + y^z + z^x) = \operatorname{ord}_2(y^z + z^x) < 3 \cdot 48^{36} \log^4 X.$$
 (44)

Hence,

$$n < 12 \cdot 48^{36} \log^4 X. \tag{45}$$

By lemma 2 (1), it follows that

$$n \log n < 12 \cdot 48^{36} \log(12 \cdot 48^{36}) \log^5 X < 1703 \cdot 48^{36} \log^5 X.$$
 (46)

We conclude that

$$X \log 2 < \log(x^y + y^z + z^x) = \log n! < n \log n < 1703 \cdot 48^{36} \log^5 X.$$

Let  $C = 1703 \cdot 48^{36}/\log 2$ , and let s = 5. Since  $\log C > 30$ , it follows, by lemma 2 (2), that

$$X < C \log^6 C < 2457 \cdot 48^{36} \cdot 148^6.$$

Hence,  $\log X < 178$ , which is a contradiction.

CASE 2.  $y < 3 \cdot 48^{36} \log^4 X$ . Let p be a prime number such that  $p \mid y$ . We first show that  $p \nmid x$ . Indeed, assume that  $p \mid x$ . Since gcd (x, y, z) = 1, it follows that  $p \nmid z$ . QWe conclude that  $p \nmid n!$ , therefore n < p. Hence,

$$n$$

In particular, n satisfies inequality (45). From Case 1 we know that  $\log X < 178$ , which is a contradiction.

Suppose now that  $p \nmid x$ . Then, by theorem vdP,

$$\operatorname{ord}_{p}(x^{y} + z^{x}) = \operatorname{ord}_{p}(1 - (-x)^{-y}z^{x}) < 48^{36} \frac{p}{\log p} \log^{4} X < < 48^{36} y \log^{4} X < 3 \cdot 48^{72} \log^{8} X.$$

$$(47)$$

We distinguish 2 cases:

CASE 2.1.  $z \ge 3 \cdot 48^{72} \log^8 X$ . In this case, by lemma 2 (1) and inequality (47),

$$\frac{n}{4(p-1)} < \operatorname{ord}_{p} n! = \operatorname{ord}_{p} (y^{z} + (x^{y} + z^{x})) =$$

$$= \operatorname{ord}_{p} (x^{y} + z^{x}) < 3 \cdot 48^{72} \log^{8} X.$$

Hence,

$$n < 12(p-1) \cdot 48^{72} \log^8 X < 12y \cdot 48^{72} \log^8 X < 36 \cdot 48^{108} \log^{12} X.$$
 (48)

From lemma 2 (1) we conclude that

$$X \log 2 < \log(x^y + y^z + z^x) = \log n! < n \log n <$$

$$< 36 \cdot 48^{108} \log(36 \cdot 48^{108}) \log^{13} X < 317 \cdot 48^{109} \log^{13} X.$$
 (49)

Let  $C = 317 \cdot 48^{109}/\log 2$ , and let s = 13. Since  $\log C > 182$ , it follows, by lemma 2 (2), that

$$X < C \log^{11} C < 458 \cdot 48^{109} \ln^{14} (458 \cdot 48^{109}) < 458 \cdot 48^{109} \cdot 429^{14}$$

Hence,  $\log X < 513$ , which is a contradiction.

CASE 2.2.  $z < 3.48^{72} \log^8 X$ . By theorem vdP, it follows that

$$\operatorname{ord}_2(z^x + (x^y + y^z)) = \operatorname{ord}_2(1 - (-x^y - y^z)z^{-X}) <$$

$$<48^{36} \frac{2}{\log 2} \log(x^y + y^z) \log^3 X < 3 \cdot 48^{36} \log(x^y + y^z) \log^3 X. \tag{50}$$

We now bound  $\log(x^y + y^z)$ . Let  $y_1 = 3 \cdot 48^{36} \log^4 X$  and  $z_1 = 3 \cdot 48^{72} \log^8 X$ . Since  $y < y_1$  and  $z < z_1$ , it follows that

$$\log(x^y + y^z) < \log(X^{y_1} + y_1^{z_1}) < \log 2 + \max(y_1 \log X, z_1 \log y_1).$$

Since  $z_1 \log y_1 > z_1 > y_1 \log X$ , it follows that

$$\log(x^y + y^z) < \log 2 + z_1 \log y_1.$$

From lemma 2 (1) we conclude that

$$\log(x^y + y^z) < \log 2 + z_1 \log y_1 = \log 2 + \frac{z_1}{y_1} \cdot (y_1 \log y_1) < \frac{z_1}{y_1} \cdot (y_2 \log y_2) < \frac{z_1}{y_2} \cdot (y_2 \log y_2) < \frac{z_1}{y_2} \cdot (y_2 \log y_2) < \frac{z_2}{y_2} \cdot (y_$$

$$< \log 2 + 48^{36} \log^4 X \cdot \left( 3 \cdot 48^{36} \log(3 \cdot 48^{36}) \right) \log^5 X < 422 \cdot 48^{72} \log^9 X.$$
 (51)

From lemma 1 and inequalities (50) and (51) it follows that

$$n/4 < \operatorname{ord}_2 n! = \operatorname{ord}_2 (z^x + (x^y + y^z)) < 1266 \cdot 48^{108} \log^{12} X.$$

Hence,

$$n < 5064 \cdot 48^{108} \log^{12} X.$$

By lemma 2 (1), it follows that

$$X \log 2 < \log(x^y + y^z + z^x) = \log n! < n \log n <$$

$$< 5064 \cdot 48^{108} \cdot \log(5064 \cdot 48^{108}) \log^{13} X < 22 \cdot 48^{111} \log^{13} X.$$

Let  $C = 22 \cdot 48^{111}/\log 2$ , and let s = 13. Since  $\log C > 182$ , it follows, by lemma 2 (2), that

$$X < C \log^{14} C < 22 \cdot 48^{111} \cdot 433^{14}$$
.

Hence,  $\log X < 518$ , which is the final contradiction.

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