

Funcoids and Reloids*

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Abstract

It is a part of my Algebraic General Topology research.

In this article I introduce the concepts of *funcoids* which generalize proximity spaces and *reloids* which generalize uniform spaces. The concept of funcoid is generalized concept of proximity space, the concept of reloid is cleared from superfluous details (generalized) concept of uniform space. Also funcoids and reloids are generalizations of binary relations whose domains and ranges are filters (instead of sets).

That funcoids and reloids are common generalizations of both (proximity, pretopology, uniform) spaces and of (multivalued) functions, makes this theory smart for analyzing properties (e.g. continuousness) of functions on spaces. Also funcoids and reloids can be considered as a generalization of (oriented) graphs, this provides us with a common generalization of analysis and discrete mathematics.

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1 Common

1.1 Draft status

This article is a draft.

This text refers to a preprint edition of [5]. Theorem number clashes may appear due editing both of these manuscripts.

1.2 Used concepts, notation and statements

Set of functions from a set A to a set B is denoted as B^A .

I will often skip parentheses and write fx instead of $f(x)$ to denote the result of a function f acting on the argument x .

I will denote $\langle f \rangle X = \{f\alpha \mid \alpha \in X\}$.

For simplicity I will assume that all sets in consideration are subsets of universal set \mathcal{U} .

1.2.1 Filters

In this work the word *filter* will refer to a filter on a set \mathcal{U} (in contrast to [5] where are considered filters on arbitrary posets).

I will call the set of filters ordered reverse to set-theoretic inclusion of filters *the set of filter objects* \mathfrak{F} and its element *filter objects* (f.o. for short). I will denote $\text{up}\mathcal{F}$ the filter corresponding to a filter object \mathcal{F} . So we have $\mathcal{A} \subseteq \mathcal{B} \Leftrightarrow \text{up}\mathcal{A} \supseteq \text{up}\mathcal{B}$ for every filter objects \mathcal{A} and \mathcal{B} . We also will equate filter objects corresponding to principal filters with corresponding sets. (Thus we have $\mathcal{P}\mathcal{U} \subseteq \mathfrak{F}$.) See [5] for formal definition of filter objects in the framework of ZF. Filters (and filter objects) are studied in the work [5].

Filter objects corresponding to ultrafilters are atoms of the lattice \mathfrak{F} and will be called *atomic filter objects*.

Also we will need to introduce the concept of *generalized filter base*.

Definition 1. *Generalized filter base* is a set $S \in \mathcal{P}\mathfrak{F} \setminus \{\emptyset\}$ such that

$$\forall \mathcal{A}, \mathcal{B} \in S \exists \mathcal{C} \in S: \mathcal{C} \subseteq \mathcal{A} \cap \mathcal{B}.$$

Proposition 2. Let S is a generalized filter base. If $A_1, \dots, A_n \in S$ ($n \in \mathbb{N}$), then

$$\exists \mathcal{C} \in S: \mathcal{C} \subseteq A_1 \cap \dots \cap A_n.$$

Proof. Can be easily proved by induction. □

Theorem 3. If S is a generalized filter base, then $\text{up} \bigcap^{\mathfrak{F}} S = \bigcup \langle \text{up} \rangle S$.

Proof. Obviously $\text{up} \bigcap^{\mathfrak{F}} S \supseteq \bigcup \langle \text{up} \rangle S$. Reversely, let $K \in \bigcap^{\mathfrak{F}} S$; then $K = A_1 \cap \dots \cap A_n$ where $A_i \in \text{up} \mathcal{A}_i \in S$, $i = 1, \dots, n$, $n \in \mathbb{N}$; so exists $\mathcal{C} \in S$ such that $\mathcal{C} \subseteq A_1 \cap \dots \cap A_n \subseteq A_1 \cap \dots \cap A_n = K$, $K \in \text{up} \mathcal{C}$, $K \in \bigcup \langle \text{up} \rangle S$. □

Corollary 4. If S is a generalized filter base, then $\bigcap^{\mathfrak{F}} S = \emptyset \Leftrightarrow \emptyset \in S$.

Proof. $\bigcap^{\mathfrak{F}} S = \emptyset \Leftrightarrow \emptyset \in \text{up} \bigcap^{\mathfrak{F}} S \Leftrightarrow \emptyset \in \bigcup \langle \text{up} \rangle S \Leftrightarrow \exists \mathcal{X} \in S: \emptyset \in \text{up} \mathcal{X} \Leftrightarrow \emptyset \in S$. □

Definition 5. I will call a *partially ordered (pre)category* a (pre)category together with partial order on each of its Hom-sets.

1.3 Earlier works

Some mathematician were researching generalizations of proximities and uniformities before me but they have failed to reach the right degree of generalization which is presented in this work allowing to represent properties of spaces with algebraic (or categorical) formulas.

Some references to predecessors:

- In [1] and [2] are studied semi-uniformities and proximities.
- [3] and [4] contains recent progress in quasi-uniform spaces.

2 Partially ordered dagger categories

2.1 Partially ordered categories

Definition 6. I will call a *partially ordered (pre)category* a (pre)category together with partial order \subseteq on each of its Hom-sets with the additional requirement that

$$f_1 \subseteq f_2 \wedge g_1 \subseteq g_2 \Rightarrow g_1 \circ f_1 \subseteq g_2 \circ f_2$$

for any morphisms f_1, g_1, f_2, g_2 such that $\text{Src } f_1 = \text{Src } f_2 \wedge \text{Dst } f_1 = \text{Dst } f_2 = \text{Src } g_1 = \text{Src } g_2 \wedge \text{Dst } g_1 = \text{Dst } g_2$.

2.2 Dagger categories

Definition 7. I will call a *dagger precategory* a precategory together with an involutive contravariant identity-on-objects prefunctor $x \mapsto x^\dagger$.

In other words, a *dagger precategory* is a precategory equipped with a function $x \mapsto x^\dagger$ on its set of morphisms which reverses the source and the destination and is subject to the following identities for every morphisms f and g :

1. $f^{\dagger\dagger} = f$;

$$2. (g \circ f)^\dagger = f^\dagger \circ g^\dagger.$$

Definition 8. I will call a *dagger category* a category together with an involutive contravariant identity-on-objects functor $x \mapsto x^\dagger$.

In other words, a *dagger category* is a category equipped with a function $x \mapsto x^\dagger$ on its set of morphisms which reverses the source and the destination and is subject to the following identities for every morphisms f and g and object A :

1. $f^{\dagger\dagger} = f$;
2. $(g \circ f)^\dagger = f^\dagger \circ g^\dagger$;
3. $(1_A)^\dagger = 1_A$.

Theorem 9. If a category is a dagger precategory then it is a dagger category.

Proof. We need to prove only that $(1_A)^\dagger = 1_A$. Really

$$(1_A)^\dagger = (1_A)^\dagger \circ 1_A = (1_A)^\dagger \circ (1_A)^{\dagger\dagger} = ((1_A)^\dagger \circ 1_A)^\dagger = (1_A)^{\dagger\dagger} = 1_A. \quad \square$$

For a partially ordered dagger (pre)category I will additionally require (for any morphisms f and g)

$$f^\dagger \subseteq g^\dagger \Leftrightarrow f \subseteq g.$$

An example of dagger category is the category **Rel** whose objects are sets and whose morphisms are binary relations between these sets with usual composition of binary relations and with $f^\dagger = f^{-1}$.

Definition 10. A morphism f of a dagger category is called *unitary* when it is an isomorphism and $f^\dagger = f^{-1}$.

Definition 11. *Symmetric* (endo)morphism of a dagger precategory is such a morphism f that $f = f^\dagger$.

Definition 12. *Transitive* (endo)morphism of a precategory is such a morphism f that $f = f \circ f$.

Theorem 13. The following conditions are equivalent for a morphism f of a dagger precategory:

1. f is symmetric and transitive.
2. $f = f^\dagger \circ f$.

Proof.

(1) \Rightarrow (2). If f is symmetric and transitive then $f^\dagger \circ f = f \circ f = f$.

(2) \Rightarrow (1). $f^\dagger = (f^\dagger \circ f)^\dagger = f^\dagger \circ f^{\dagger\dagger} = f^\dagger \circ f = f$, so f is symmetric. $f = f^\dagger \circ f = f \circ f$, so f is transitive. \square

2.2.1 Monovalued and entirely defined morphisms

Definition 14. For a partially ordered dagger category I will call *monovalued* morphism such a morphism f that $f \circ f^\dagger \subseteq 1_{\text{Dst } f}$.

Definition 15. For a partially ordered dagger category I will call *entirely defined* morphism such a morphism f that $f^\dagger \circ f \supseteq 1_{\text{Src } f}$.

Remark 16. Easy to show that this is a generalization of monovalued and entirely defined binary relations as morphisms of the category **Rel**.

Definition 17. For a given partially ordered dagger category C the *category of monovalued (entirely defined) morphisms* of C is the category with the same set of objects as of C and the set of morphisms being the set of monovalued (entirely defined) morphisms of C with the composition of morphisms the same as in C .

We need to prove that these are really categories, that is that composition of monovalued (entirely defined) morphisms is monovalued (entirely defined) and that identity morphisms are monovalued and entirely defined.

Proof.

Monovalued. Let f and g are monovalued morphisms, $\text{Dst } f = \text{Src } g$. $(g \circ f) \circ (g \circ f)^\dagger = g \circ f \circ f^\dagger \circ g^\dagger \subseteq g \circ 1_{\text{Dst } f} \circ g^\dagger = g \circ 1_{\text{Src } g} \circ g^\dagger = g \circ g^\dagger \subseteq 1_{\text{Dst } g} = 1_{\text{Dst}(g \circ f)}$. So $g \circ f$ is monovalued.

That identity morphisms are monovalued follows from the following: $1_A \circ (1_A)^\dagger = 1_A \circ 1_A = 1_A = 1_{\text{Dst } 1_A} \subseteq 1_{\text{Dst } 1_A}$.

Entirely defined. Let f and g are entirely defined morphisms, $\text{Dst } f = \text{Src } g$. $(g \circ f)^\dagger \circ (g \circ f) = f^\dagger \circ g^\dagger \circ g \circ f \supseteq f^\dagger \circ 1_{\text{Src } g} \circ f = f^\dagger \circ 1_{\text{Dst } f} \circ f = f^\dagger \circ f \supseteq 1_{\text{Src } f} = 1_{\text{Src}(g \circ f)}$. So $g \circ f$ is entirely defined.

That identity morphisms are entirely defined follows from the following: $(1_A)^\dagger \circ 1_A = 1_A \circ 1_A = 1_A = 1_{\text{Src } 1_A} \subseteq 1_{\text{Src } 1_A}$. \square

3 Funcoids

3.1 Informal introduction into funcoids

Funcoids are a generalization of proximity spaces and a generalization of pretopological spaces. Also funcoids are a generalization of binary relations.

That funcoids are a common generalization of “spaces” (proximity spaces, (pre)topological spaces) and binary relations (including monovalued functions) makes them smart for describing properties of functions in regard of spaces. For example the statement “ f is a continuous function from a space μ to a space ν ” can be described in terms of funcoids as the formula $f \circ \mu \subseteq \nu \circ f$ (see my yet unpublished article “Generalized continuity” for details).

Most naturally funcoids appear as a generalization of proximity spaces.

Let δ be a proximity that is certain binary relation so that $A \delta B$ is defined for any sets A and B . We will extend it from sets to filter objects by the formula:

$$\mathcal{A} \delta' \mathcal{B} \Leftrightarrow \forall A \in \text{up } \mathcal{A}, B \in \text{up } \mathcal{B}: A \delta B.$$

Then (as will be proved below) exist two functions $\alpha, \beta \in \mathfrak{F}^\delta$ such that

$$\mathcal{A} \delta' \mathcal{B} \Leftrightarrow \mathcal{B} \cap \alpha \mathcal{A} \neq \emptyset \Leftrightarrow \mathcal{A} \cap \beta \mathcal{B} \neq \emptyset.$$

The pair $(\alpha; \beta)$ is called *funcoid* when $\mathcal{B} \cap \alpha \mathcal{A} \neq \emptyset \Leftrightarrow \mathcal{A} \cap \beta \mathcal{B} \neq \emptyset$. So funcoids are a generalization of proximity spaces.

Funcoids consist of two components the first α and the second β . The first component of a funcoid f is denoted as $\langle f \rangle$ and the second component is denoted as $\langle f^{-1} \rangle$. (The similarity of this notation with the notation for the image of a set under a function is not a coincidence, we will see that in the case of discrete funcoids (see below) these coincide.)

One of the most important properties of a funcoid is that it is uniquely determined by just one of its components. That is a funcoid f is uniquely determined by the function $\langle f \rangle$. Moreover a funcoid f is uniquely determined by $\langle f \rangle|_{\mathcal{P}U}$ that is by values of function $\langle f \rangle$ on sets.

Next we will consider some examples of funcoids determined by specified values of the first component on sets.

Funcoids as a generalization of pretopological spaces: Let α be a pretopological space that is a map $\alpha \in \mathfrak{F}^{\cup}$. Then we define $\alpha'X \stackrel{\text{def}}{=} \bigcup^{\mathfrak{F}} \{\alpha X \mid x \in X\}$ for any set X . We will prove that there exists a unique funcoid f such that $\alpha' = \langle f \rangle|_{\mathcal{P}U}$. So funcoids are a generalization of pretopological spaces. Funcoids are also a generalization of preclosure operators: For every preclosure operator p exists unique funcoid such that $\langle f \rangle|_{\mathcal{P}U} = p$; in this case $\langle f \rangle|_{\mathcal{P}U} \in \mathcal{P}U^{\mathcal{P}U}$.

For any binary relation p exists unique funcoid f such that $\forall X \in \mathcal{P}U: \langle f \rangle X = \langle p \rangle X$ (where $\langle p \rangle$ is defined in the introduction), recall that a funcoid is uniquely determined by the values of its first component on sets. I will call such funcoids *discrete*. So funcoids are a generalization of binary relations.

Composition of binary relations (i.e. of discrete funcoids) complies with the formulas:

$$\langle g \circ f \rangle = \langle g \rangle \circ \langle f \rangle \quad \text{and} \quad \langle (g \circ f)^{-1} \rangle = \langle f^{-1} \rangle \circ \langle g^{-1} \rangle.$$

By the same formulas we can define composition of any two funcoids.

Also funcoids can be reversed (like reversal of X and Y in a binary relation) by the formula $(\alpha; \beta)^{-1} = (\beta; \alpha)$. In particular case if μ is a proximity we have $\mu^{-1} = \mu$ because proximities are symmetric.

Funcoids behave similarly to (multivalued) functions but acting on filter objects instead of acting on sets. Below will be defined domain and image of a funcoid (the domain and the image of a funcoid are filter objects).

3.2 Basic definitions

Definition 18. Let's call a *funcoid* a pair $(\alpha; \beta)$ where $\alpha, \beta \in \mathfrak{F}^{\mathfrak{F}}$ such that

$$\forall \mathcal{X}, \mathcal{Y} \in \mathfrak{F}: (\mathcal{Y} \cap^{\mathfrak{F}} \alpha \mathcal{X} \neq \emptyset \Leftrightarrow \mathcal{X} \cap^{\mathfrak{F}} \beta \mathcal{Y} \neq \emptyset).$$

Definition 19. $\langle (\alpha; \beta) \rangle \stackrel{\text{def}}{=} \alpha$ for a funcoid $(\alpha; \beta)$.

Definition 20. $(\alpha; \beta)^{-1} = (\beta; \alpha)$ for a funcoid $(\alpha; \beta)$.

Proposition 21. If f is a funcoid then f^{-1} is also a funcoid.

Proof. Follows from symmetry in the definition of funcoid. □

Obvious 22. $(f^{-1})^{-1} = f$ for a funcoid f .

Definition 23. The relation $[f] \in \mathcal{P}\mathfrak{F}^2$ is defined by the formula (for any filter objects \mathcal{X}, \mathcal{Y} and funcoid f)

$$\mathcal{X}[f]\mathcal{Y} \stackrel{\text{def}}{=} \mathcal{Y} \cap^{\mathfrak{F}} \langle f \rangle \mathcal{X} \neq \emptyset.$$

Obvious 24. $\mathcal{X}[f]\mathcal{Y} \Leftrightarrow \mathcal{Y} \cap^{\mathfrak{F}} \langle f \rangle \mathcal{X} \neq \emptyset \Leftrightarrow \mathcal{X} \cap^{\mathfrak{F}} \langle f^{-1} \rangle \mathcal{Y}$ for any filter objects \mathcal{X}, \mathcal{Y} and funcoid f .

Obvious 25. $[f^{-1}] = [f]^{-1}$ for a funcoid f .

Theorem 26.

1. For given value of $\langle f \rangle$ exists no more than one funcoid f .
2. For given value of $[f]$ exists no more than one funcoid f .

Proof. Let f and g are funcoids.

Obviously $\langle f \rangle = \langle g \rangle \Rightarrow [f] = [g]$ and $\langle f^{-1} \rangle = \langle g^{-1} \rangle \Rightarrow [f] = [g]$. So enough to prove that $[f] = [g] \Rightarrow \langle f \rangle = \langle g \rangle$.

Provided that $[f] = [g]$ we have $\mathcal{Y} \cap^{\mathfrak{F}} \langle f \rangle \mathcal{X} \neq \emptyset \Leftrightarrow \mathcal{X}[f]\mathcal{Y} \Leftrightarrow \mathcal{X}[g]\mathcal{Y} \Leftrightarrow \mathcal{Y} \cap^{\mathfrak{F}} \langle g \rangle \mathcal{X} \neq \emptyset$ and consequently $\langle f \rangle \mathcal{X} = \langle g \rangle \mathcal{X}$ for any f.o. \mathcal{X} and \mathcal{Y} because the set of filter objects is separable [5], thus $\langle f \rangle = \langle g \rangle$. □

Proposition 27. $\langle f \rangle(\mathcal{I} \cup^{\mathfrak{F}} \mathcal{J}) = \langle f \rangle \mathcal{I} \cup^{\mathfrak{F}} \langle f \rangle \mathcal{J}$ for any funcooid f and $\mathcal{I}, \mathcal{J} \in \mathfrak{F}$.

Proof. [TODO: Point used theorems.]

$$\begin{aligned}
\star \langle f \rangle(\mathcal{I} \cup^{\mathfrak{F}} \mathcal{J}) &= \\
\{\mathcal{Y} \in \mathfrak{F} \mid \mathcal{Y} \cap^{\mathfrak{F}} \langle f \rangle(\mathcal{I} \cup^{\mathfrak{F}} \mathcal{J}) \neq \emptyset\} &= \text{(by corollary 10 in [5])} \\
\{\mathcal{Y} \in \mathfrak{F} \mid (\mathcal{I} \cup^{\mathfrak{F}} \mathcal{J}) \cap^{\mathfrak{F}} \langle f^{-1} \rangle \mathcal{Y} \neq \emptyset\} &= \\
\{\mathcal{Y} \in \mathfrak{F} \mid (\mathcal{I} \cap^{\mathfrak{F}} \langle f^{-1} \rangle \mathcal{Y}) \cup^{\mathfrak{F}} (\mathcal{J} \cap^{\mathfrak{F}} \langle f^{-1} \rangle \mathcal{Y}) \neq \emptyset\} &= \\
\{\mathcal{Y} \in \mathfrak{F} \mid (\mathcal{Y} \cap^{\mathfrak{F}} \langle f \rangle \mathcal{I}) \cup^{\mathfrak{F}} (\mathcal{Y} \cap^{\mathfrak{F}} \langle f \rangle \mathcal{J}) \neq \emptyset\} &= \\
\{\mathcal{Y} \in \mathfrak{F} \mid \mathcal{Y} \cap^{\mathfrak{F}} (\langle f \rangle \mathcal{I} \cup^{\mathfrak{F}} \langle f \rangle \mathcal{J}) \neq \emptyset\} &= \\
\star \langle \langle f \rangle \mathcal{I} \cup^{\mathfrak{F}} \langle f \rangle \mathcal{J} \rangle. &
\end{aligned}$$

Thus $\langle f \rangle(\mathcal{I} \cup^{\mathfrak{F}} \mathcal{J}) = \langle f \rangle \mathcal{I} \cup^{\mathfrak{F}} \langle f \rangle \mathcal{J}$ because \mathfrak{F} is separable. \square

3.2.1 Composition of funcooids

Definition 28. *Composition* of funcooids is defined by the formula

$$(\alpha_2; \beta_2) \circ (\alpha_1; \beta_1) = (\alpha_2 \circ \alpha_1; \beta_1 \circ \beta_2).$$

Proposition 29. If f, g are funcooids then $g \circ f$ is funcooid.

Proof. Let $f = (\alpha_1; \beta_1)$, $g = (\alpha_2; \beta_2)$.

$$\mathcal{Y} \cap^{\mathfrak{F}} (\alpha_2 \circ \alpha_1) \mathcal{X} \neq \emptyset \Leftrightarrow \mathcal{Y} \cap^{\mathfrak{F}} \alpha_2 \alpha_1 \mathcal{X} \neq \emptyset \Leftrightarrow \alpha_1 \mathcal{X} \cap^{\mathfrak{F}} \beta_2 \mathcal{Y} \neq \emptyset \Leftrightarrow \mathcal{X} \cap^{\mathfrak{F}} \beta_1 \beta_2 \mathcal{Y} \neq \emptyset \Leftrightarrow \mathcal{X} \cap^{\mathfrak{F}} (\beta_1 \circ \beta_2) \mathcal{Y} \neq \emptyset.$$

So $(\alpha_2 \circ \alpha_1; \beta_1 \circ \beta_2)$ is a funcooid. \square

Obvious 30. $\langle g \circ f \rangle = \langle g \rangle \circ \langle f \rangle$ for any funcooids f and g .

Proposition 31. $(h \circ g) \circ f = h \circ (g \circ f)$ for any funcooids f, g, h .

Proof.

$$\langle (h \circ g) \circ f \rangle = \langle h \circ g \rangle \circ \langle f \rangle = (\langle h \rangle \circ \langle g \rangle) \circ \langle f \rangle = \langle h \rangle \circ (\langle g \rangle \circ \langle f \rangle) = \langle h \rangle \circ \langle g \circ f \rangle = \langle h \circ (g \circ f) \rangle. \quad \square$$

Theorem 32. $(g \circ f)^{-1} = f^{-1} \circ g^{-1}$ for any funcooids f and g .

Proof. $\langle (g \circ f)^{-1} \rangle = \langle f^{-1} \rangle \circ \langle g^{-1} \rangle = \langle f^{-1} \circ g^{-1} \rangle. \quad \square$

3.3 Funcooid as continuation

Theorem 33. For any funcooid f and filter objects \mathcal{X} and \mathcal{Y}

1. $\langle f \rangle \mathcal{X} = \bigcap^{\mathfrak{F}} \langle \langle f \rangle \rangle \text{up } \mathcal{X}$;
2. $\mathcal{X}[f] \mathcal{Y} \Leftrightarrow \forall X \in \text{up } \mathcal{X}, Y \in \text{up } \mathcal{Y}: X[f] Y$.

Proof. 2. $\mathcal{X}[f] \mathcal{Y} \Leftrightarrow \mathcal{Y} \cap^{\mathfrak{F}} \langle f \rangle \mathcal{X} \neq \emptyset \Leftrightarrow \forall Y \in \text{up } \mathcal{Y}: Y \cap^{\mathfrak{F}} \langle f \rangle \mathcal{X} \neq \emptyset \Leftrightarrow \forall Y \in \text{up } \mathcal{Y}: \mathcal{X}[f] Y$.

Analogously $\mathcal{X}[f] \mathcal{Y} \Leftrightarrow \forall X \in \text{up } \mathcal{X}: X[f] \mathcal{Y}$. Combining these two equalities we get

$$\mathcal{X}[f] \mathcal{Y} \Leftrightarrow \forall X \in \text{up } \mathcal{X}, Y \in \text{up } \mathcal{Y}: X[f] Y.$$

1. $\mathcal{Y} \cap^{\mathfrak{F}} \langle f \rangle \mathcal{X} \neq \emptyset \Leftrightarrow \forall X \in \text{up } \mathcal{X}: \mathcal{Y} \cap^{\mathfrak{F}} \langle f \rangle X \neq \emptyset$

Let's denote $W = \{\mathcal{Y} \cap^{\mathfrak{F}} \langle f \rangle X \mid X \in \text{up } \mathcal{X}\}$. Let $\mathcal{P}, \mathcal{Q} \in W$. Then $\mathcal{P} = \mathcal{Y} \cap^{\mathfrak{F}} \langle f \rangle A$, $\mathcal{Q} = \mathcal{Y} \cap^{\mathfrak{F}} \langle f \rangle B$ where $A, B \in \text{up } \mathcal{X}$; $A \cap B \in \text{up } \mathcal{X}$ and $\mathcal{R} \subseteq \mathcal{P} \cap^{\mathfrak{F}} \mathcal{Q}$ for $\mathcal{R} = \mathcal{Y} \cap^{\mathfrak{F}} \langle f \rangle (A \cap B) \in W$. So W is a generalized filter base. [TODO: Simplify the proof than it is a g.f.b.]

$\emptyset \notin W \Leftrightarrow \bigcap^{\mathfrak{F}} W \neq \emptyset$ by the corollary 4 of the theorem 3. That is

$$\forall X \in \text{up } \mathcal{X}: \mathcal{Y} \cap^{\mathfrak{F}} \langle f \rangle X \neq \emptyset \Leftrightarrow \mathcal{Y} \cap^{\mathfrak{F}} \bigcap^{\mathfrak{F}} \langle \langle f \rangle \rangle \text{up } \mathcal{X} \neq \emptyset.$$

Comparing with the above, $\mathcal{Y} \cap^{\mathfrak{F}} \langle f \rangle \mathcal{X} \neq \emptyset \Leftrightarrow \mathcal{Y} \cap^{\mathfrak{F}} \bigcap^{\mathfrak{F}} \langle \langle f \rangle \rangle \text{up } \mathcal{X} \neq \emptyset$. So $\langle f \rangle \mathcal{X} = \bigcap^{\mathfrak{F}} \langle \langle f \rangle \rangle \text{up } \mathcal{X}$. \square

Theorem 34.

1. A function $\alpha \in \mathfrak{F}^{\mathcal{P}\mathcal{U}}$ conforming to the formulas (for any $I, J \in \mathcal{P}\mathcal{U}$)

$$\alpha \emptyset = \emptyset, \quad \alpha(I \cup J) = \alpha I \cup \alpha J$$

can be continued to the function $\langle f \rangle$ for exactly one funcoid f ;

$$\langle f \rangle \mathcal{X} = \bigcap^{\mathfrak{F}} \langle \alpha \rangle \text{up } \mathcal{X} \tag{1}$$

for any filter object \mathcal{X} .

2. A relation $\delta \in \mathcal{P}(\mathcal{P}\mathcal{U})^2$ conforming to the formulas (for any $I, J, K \in \mathcal{P}\mathcal{U}$)

$$\begin{aligned} \neg(\emptyset \delta I), \quad I \cup J \delta K &\Leftrightarrow I \delta K \vee J \delta K, \\ \neg(I \delta \emptyset), \quad K \delta I \cup J &\Leftrightarrow K \delta I \vee K \delta J \end{aligned} \tag{2}$$

can be continued to the relation $[f]$ for exactly one funcoid f ;

$$\mathcal{X}[f]\mathcal{Y} \Leftrightarrow \forall X \in \text{up } \mathcal{X}, Y \in \text{up } \mathcal{Y}: X \delta Y \tag{3}$$

for any filter objects \mathcal{X}, \mathcal{Y} .

Proof. Existence of no more than one such funcoids and formulas (1) and (3) follow from the previous theorem.

2. Let define $\alpha \in \mathfrak{F}^{\mathcal{P}\mathcal{U}}$ by the formula $\partial(\alpha X) = \{Y \in \mathcal{P}\mathcal{U} \mid X \delta Y\}$ for any $X \in \mathcal{P}\mathcal{U}$. Analogously can be defined $\beta \in \mathfrak{F}^{\mathcal{P}\mathcal{U}}$ by the formula $\partial(\beta X) = \{X \in \mathcal{P}\mathcal{U} \mid X \delta Y\}$. Let's continue α and β to $\alpha' \in \mathfrak{F}^{\mathfrak{F}}$ and $\beta' \in \mathfrak{F}^{\mathfrak{F}}$ by the formulas

$$\alpha' \mathcal{X} = \bigcap^{\mathfrak{F}} \langle \alpha \rangle \text{up } \mathcal{X} \quad \text{and} \quad \beta' \mathcal{X} = \bigcap^{\mathfrak{F}} \langle \beta \rangle \text{up } \mathcal{X}.$$

and δ to $\delta' \in \mathcal{P}\mathfrak{F}^2$ by the formula

$$\mathcal{X} \delta' \mathcal{Y} \Leftrightarrow \forall X \in \text{up } \mathcal{X}, Y \in \text{up } \mathcal{Y}: X \delta Y.$$

$\mathcal{Y} \cap^{\mathfrak{F}} \alpha' \mathcal{X} \neq \emptyset \Leftrightarrow \mathcal{Y} \cap^{\mathfrak{F}} \bigcap^{\mathfrak{F}} \langle \alpha \rangle \text{up } \mathcal{X} \neq \emptyset \Leftrightarrow \bigcap^{\mathfrak{F}} \langle \mathcal{Y} \cap^{\mathfrak{F}} \rangle \langle \alpha \rangle \text{up } \mathcal{X} \neq \emptyset$. Let's prove that

$$W = \langle \mathcal{Y} \cap^{\mathfrak{F}} \rangle \langle \alpha \rangle \text{up } \mathcal{X}$$

is a generalized filter base: **[TODO: Simplify this proof.]** If $\mathcal{A}, \mathcal{B} \in W$ then exist $X_1, X_2 \in \text{up } \mathcal{X}$ such that

$$\mathcal{A} = \mathcal{Y} \cap^{\mathfrak{F}} \alpha X_1 \quad \text{and} \quad \mathcal{B} = \mathcal{Y} \cap^{\mathfrak{F}} \alpha X_2.$$

Then $\mathcal{Y} \cap^{\mathfrak{F}} \alpha(X_1 \cap X_2) \in W$. So W is a generalized filter base.

Accordingly the corollary 4 of the theorem 3, $\bigcap^{\mathfrak{F}} \langle \mathcal{Y} \cap^{\mathfrak{F}} \rangle \langle \alpha \rangle \text{up } \mathcal{X} \neq \emptyset$ is equivalent to

$$\forall X \in \text{up } \mathcal{X}: \mathcal{Y} \cap^{\mathfrak{F}} \alpha X \neq \emptyset,$$

what is equivalent to $\forall X \in \text{up } \mathcal{X}, Y \in \text{up } \mathcal{Y}: Y \cap^{\mathfrak{F}} \alpha X \neq \emptyset \Leftrightarrow \forall X \in \text{up } \mathcal{X}, Y \in \text{up } \mathcal{Y}: Y \in \partial(\alpha X) \Leftrightarrow \forall X \in \text{up } \mathcal{X}, Y \in \text{up } \mathcal{Y}: X \delta Y$. Combining the equivalencies we get $\mathcal{Y} \cap^{\mathfrak{F}} \alpha' \mathcal{X} \neq \emptyset \Leftrightarrow X \delta' Y$. Analogously $\mathcal{X} \cap^{\mathfrak{F}} \beta' \mathcal{Y} \neq \emptyset \Leftrightarrow X \delta' Y$. So $\mathcal{Y} \cap^{\mathfrak{F}} \alpha' \mathcal{X} \neq \emptyset \Leftrightarrow \mathcal{X} \cap^{\mathfrak{F}} \beta' \mathcal{Y} \neq \emptyset$, that is $(\alpha'; \beta')$ is a funcoid. From the formula $\mathcal{Y} \cap^{\mathfrak{F}} \alpha' \mathcal{X} \neq \emptyset \Leftrightarrow X \delta' Y$ follows that $[(\alpha'; \beta')]$ is a continuation of δ .

1. Let define the relation $\delta \in \mathcal{P}(\mathcal{P}\mathcal{U})^2$ by the formula $X \delta Y \Leftrightarrow Y \cap^{\mathfrak{F}} \alpha X \neq \emptyset$. Then the formulas (2) are true.

Accordingly the above δ can be continued to the relation $[f]$ for some funcoid f .

$\forall X, Y \in \mathcal{P}\mathcal{U}: X[f]Y \Leftrightarrow Y \cap^{\mathfrak{F}} \alpha X \neq \emptyset$, consequently $\forall X \in \mathcal{P}\mathcal{U}: \alpha X = \langle f \rangle X$. So $\langle f \rangle$ is a continuation of α . \square

Note that by the last theorem to every proximity δ corresponds exactly one funcoid. So funcoids is a generalization of proximity structures.

Definition 35. Any (multivalued) function f will be considered as a funcoid, where by definition $\langle f \rangle \mathcal{X} = \bigcap^{\delta} \langle \langle f \rangle \rangle \text{up} \mathcal{X}$ for any $\mathcal{X} \in \mathfrak{F}$.

Using the last theorem it is easy to show that this definition is monovalued and does not contradict to former stuff.

So binary relations (= multivalued functions) may be called *discrete funcoids*.

I will denote FCD the set of funcoids or the category of funcoids (see below) dependently on context.

3.4 Lattice of funcoids

Definition 36. $f \subseteq g \stackrel{\text{def}}{=} [f] \subseteq [g]$ for $f, g \in \text{FCD}$.

Thus FCD is a poset.

Conjecture 37. The filtrator of funcoids is:

1. with separable core;
2. with co-separable core.

Definition 38. I will call the *filtrator of funcoids* (see [5] for the definition of filtrators) the filtrator (FCD; $\mathcal{P}\mathcal{U}^2$).

Theorem 39. The set of funcoids is a complete lattice. For any $R \in \mathcal{P}\text{FCD}$ and $X, Y \in \mathcal{P}\mathcal{U}$

1. $X[\bigcup^{\text{FCD}} R]Y \Leftrightarrow \exists f \in R: X[f]Y$;
2. $\langle \bigcup^{\text{FCD}} R \rangle X = \bigcup^{\delta} \{ \langle f \rangle X \mid f \in R \}$.

Proof.

2. $\langle h \rangle X \stackrel{\text{def}}{=} \bigcup^{\delta} \{ \langle f \rangle X \mid f \in R \}$. $\langle h \rangle \emptyset = \emptyset$;

$$\begin{aligned} \langle h \rangle (I \cup J) &= \bigcup^{\delta} \{ \langle f \rangle (I \cup J) \mid f \in R \} \\ &= \bigcup^{\delta} \{ \langle f \rangle (I \cup^{\delta} J) \mid f \in R \} \\ &= \bigcup^{\delta} \{ \langle f \rangle I \cup^{\delta} \langle f \rangle J \mid f \in R \} \\ &= \bigcup^{\delta} \{ \langle f \rangle I \mid f \in R \} \cup^{\delta} \bigcup^{\delta} \{ \langle f \rangle J \mid f \in R \} \\ &= \langle h \rangle I \cup^{\delta} \langle h \rangle J. \end{aligned}$$

So $\langle h \rangle$ can be continued to a funcoid. Obviously

$$\forall f \in R: h \supseteq f. \quad (4)$$

And h is the least funcoid for which holds the condition (4). So $h = \bigcup^{\text{FCD}} R$.

1. $X[\bigcup^{\text{FCD}} R]Y \Leftrightarrow Y \cap^{\delta} \langle \bigcup^{\text{FCD}} R \rangle X \neq \emptyset \Leftrightarrow Y \cap^{\delta} \bigcup^{\delta} \{ \langle f \rangle X \mid f \in R \} \neq \emptyset \Leftrightarrow \exists f \in R: Y \cap^{\delta} \langle f \rangle X \neq \emptyset \Leftrightarrow \exists f \in R: X[f]Y$. \square

In the next theorem, compared to the previous one, the class of infinite unions is replaced with lesser class of finite unions and simultaneously class of sets is changed to more wide class of filter objects.

Theorem 40. For any funcoids f and g and a filter object \mathcal{X}

1. $\langle f \cup^{\text{FCD}} g \rangle \mathcal{X} = \langle f \rangle \mathcal{X} \cup^{\delta} \langle g \rangle \mathcal{X}$;
2. $[f \cup^{\delta} g] = [f] \cup [g]$.

Proof.

1. Let $\langle h \rangle \mathcal{X} \stackrel{\text{def}}{=} \langle f \rangle \mathcal{X} \cup^{\mathfrak{F}} \langle g \rangle \mathcal{X}$; $\langle h^{-1} \rangle \mathcal{Y} \stackrel{\text{def}}{=} \langle f^{-1} \rangle \mathcal{Y} \cup^{\mathfrak{F}} \langle g^{-1} \rangle \mathcal{Y}$ for any $\mathcal{X}, \mathcal{Y} \in \mathfrak{F}$. Then

$$\begin{aligned} \mathcal{Y} \cap^{\mathfrak{F}} \langle h \rangle \mathcal{X} \neq \emptyset &\Leftrightarrow \mathcal{Y} \cap^{\mathfrak{F}} \langle f \rangle \mathcal{X} \neq \emptyset \vee \mathcal{Y} \cap^{\mathfrak{F}} \langle g \rangle \mathcal{X} \neq \emptyset \\ &\Leftrightarrow \mathcal{X} \cap^{\mathfrak{F}} \langle f^{-1} \rangle \mathcal{Y} \neq \emptyset \vee \mathcal{X} \cap^{\mathfrak{F}} \langle g^{-1} \rangle \mathcal{Y} \neq \emptyset \\ &\Leftrightarrow \mathcal{X} \cap^{\mathfrak{F}} \langle h^{-1} \rangle \mathcal{Y} \neq \emptyset. \end{aligned}$$

So h is a funcoid. Consequently $f \cup^{\text{FCD}} g = h$.

2. $\mathcal{X}[f \cup^{\text{FCD}} g] \mathcal{Y} \Leftrightarrow \mathcal{Y} \cap^{\mathfrak{F}} \langle f \cup^{\text{FCD}} g \rangle \mathcal{X} \neq \emptyset \Leftrightarrow \mathcal{Y} \cap^{\mathfrak{F}} (\langle f \rangle \mathcal{X} \cup^{\mathfrak{F}} \langle g \rangle \mathcal{X}) \neq \emptyset \Leftrightarrow \mathcal{Y} \cap^{\mathfrak{F}} \langle f \rangle \mathcal{X} \neq \emptyset \vee \mathcal{Y} \cap^{\mathfrak{F}} \langle g \rangle \mathcal{X} \neq \emptyset \Leftrightarrow \mathcal{X}[f] \mathcal{Y} \vee \mathcal{X}[g] \mathcal{Y}$ for any $\mathcal{X}, \mathcal{Y} \in \mathfrak{F}$. \square

3.5 More on composition of funcoids

Proposition 41. $[g \circ f] = [g] \circ \langle f \rangle = \langle g^{-1} \rangle^{-1} \circ [f]$ for $f, g \in \text{FCD}$.

Proof. $\mathcal{X}[g \circ f] \mathcal{Y} \Leftrightarrow \mathcal{Y} \cap^{\mathfrak{F}} \langle g \circ f \rangle \mathcal{X} \neq \emptyset \Leftrightarrow \mathcal{Y} \cap^{\mathfrak{F}} \langle g \rangle \langle f \rangle \mathcal{X} \neq \emptyset \Leftrightarrow \langle f \rangle \mathcal{X}[g] \mathcal{Y} \Leftrightarrow \mathcal{X}([g] \circ \langle f \rangle) \mathcal{Y}$ for any $\mathcal{X}, \mathcal{Y} \in \mathfrak{F}$. $[g \circ f] = [(f^{-1} \circ g^{-1})^{-1}] = [f^{-1} \circ g^{-1}]^{-1} = ([f^{-1}] \circ \langle g^{-1} \rangle)^{-1} = \langle g^{-1} \rangle^{-1} \circ [f]$. \square

The following theorem is the variant for funcoids of the statement (which defines compositions of relations) that $x(g \circ f)z \Leftrightarrow \exists y(xfy \wedge ygz)$ for any x and z and any binary relations f and g .

Theorem 42. For any $\mathcal{X}, \mathcal{Z} \in \mathfrak{F}$ and $f, g \in \text{FCD}$

$$\mathcal{X}[g \circ f] \mathcal{Z} \Leftrightarrow \exists y \in \text{atoms}^{\mathfrak{F}} \mathcal{U}: (\mathcal{X}[f]y \wedge y[g] \mathcal{Z}).$$

Proof.

$$\begin{aligned} \exists y \in \text{atoms}^{\mathfrak{F}} \mathcal{U}: (\mathcal{X}[f]y \wedge y[g] \mathcal{Z}) &\Leftrightarrow \exists y \in \text{atoms}^{\mathfrak{F}} \mathcal{U}: (\mathcal{Z} \cap^{\mathfrak{F}} \langle g \rangle y \neq \emptyset \wedge y \cap^{\mathfrak{F}} \langle f \rangle \mathcal{X} \neq \emptyset) \\ &\Leftrightarrow \exists y \in \text{atoms}^{\mathfrak{F}} \mathcal{U}: (\mathcal{Z} \cap^{\mathfrak{F}} \langle g \rangle y \neq \emptyset \wedge y \subseteq \langle f \rangle \mathcal{X}) \\ &\Rightarrow \mathcal{Z} \cap^{\mathfrak{F}} \langle g \rangle \langle f \rangle \mathcal{X} \neq \emptyset \\ &\Leftrightarrow \mathcal{X}[g \circ f] \mathcal{Z}. \end{aligned}$$

Reversely, if $\mathcal{X}[g \circ f] \mathcal{Z}$ then $\langle f \rangle \mathcal{X}[g] \mathcal{Z}$, consequently exists $y \in \text{atoms}^{\mathfrak{F}} \langle f \rangle \mathcal{X}$ such that $y[g] \mathcal{Z}$; we have $\mathcal{X}[f]y$. \square

Theorem 43. If f, g, h are funcoids then

1. $f \circ (g \cup^{\text{FCD}} h) = f \circ g \cup^{\text{FCD}} f \circ h$;
2. $(g \cup^{\text{FCD}} h) \circ f = g \circ f \cup^{\text{FCD}} h \circ f$.

Proof. I will prove only the first equality because the other is analogous.

For any $\mathcal{X}, \mathcal{Z} \in \mathfrak{F}$

$$\begin{aligned} \mathcal{X}[f \circ (g \cup^{\text{FCD}} h)] \mathcal{Z} &\Leftrightarrow \exists y \in \text{atoms}^{\mathfrak{F}} \mathcal{U}: (\mathcal{X}[g \cup^{\text{FCD}} h]y \wedge y[f] \mathcal{Z}) \\ &\Leftrightarrow \exists y \in \text{atoms}^{\mathfrak{F}} \mathcal{U}: ((\mathcal{X}[g]y \vee \mathcal{X}[h]y) \wedge y[f] \mathcal{Z}) \\ &\Leftrightarrow \exists y \in \text{atoms}^{\mathfrak{F}} \mathcal{U}: (\mathcal{X}[g]y \wedge y[f] \mathcal{Z} \vee \mathcal{X}[h]y \wedge y[f] \mathcal{Z}) \\ &\Leftrightarrow \exists y \in \text{atoms}^{\mathfrak{F}} \mathcal{U}: (\mathcal{X}[g]y \wedge y[f] \mathcal{Z}) \vee \exists y \in \text{atoms}^{\mathfrak{F}} \mathcal{U}: (\mathcal{X}[h]y \wedge y[f] \mathcal{Z}) \\ &\Leftrightarrow \mathcal{X}[f \circ g] \mathcal{Z} \vee \mathcal{X}[f \circ h] \mathcal{Z} \\ &\Leftrightarrow \mathcal{X}[f \circ g \cup^{\text{FCD}} f \circ h] \mathcal{Z}. \end{aligned}$$

\square

3.6 Domain and range of a funcoid

Definition 44. Let $\mathcal{A} \in \mathfrak{F}$. The identity funcoid $I_{\mathcal{A}}$ will be defined by the formula $\langle I_{\mathcal{A}} \rangle \mathcal{X} \stackrel{\text{def}}{=} \mathcal{X} \cap^{\mathfrak{F}} \mathcal{A}$.

Proposition 45. This definition is correct and $\mathcal{X}[I_{\mathcal{A}}]\mathcal{Y} \Leftrightarrow \mathcal{A} \cap^{\mathfrak{F}} \mathcal{X} \cap^{\mathfrak{F}} \mathcal{Y} \neq \emptyset$ for any $\mathcal{X}, \mathcal{Y} \in \mathfrak{F}$; $(I_{\mathcal{A}})^{-1} = I_{\mathcal{A}}$.

Proof. $\mathcal{Y} \cap^{\mathfrak{F}} \langle I_{\mathcal{A}} \rangle \mathcal{X} \neq \emptyset \Leftrightarrow \mathcal{Y} \cap^{\mathfrak{F}} \mathcal{A} \cap^{\mathfrak{F}} \mathcal{X} \neq \emptyset \Leftrightarrow \mathcal{X} \cap^{\mathfrak{F}} \langle I_{\mathcal{A}} \rangle \mathcal{Y} \neq \emptyset$. \square

Definition 46. I will define *restricting* of a funcoid f to a filter object \mathcal{A} by the formula $f|_{\mathcal{A}} \stackrel{\text{def}}{=} f \circ I_{\mathcal{A}}$.

Obviously the last definition does not contradict to the previous.

Definition 47. *Image* of a funcoid f will be defined by the formula $\text{im } f = \langle f \rangle \mathfrak{U}$.

Domain of a funcoid f is defined by the formula $\text{dom } f = \text{im } f^{-1}$.

Proposition 48. $\langle f \rangle \mathcal{X} = \langle f \rangle (\mathcal{X} \cap^{\mathfrak{F}} \text{dom } f)$ for any $f \in \text{FCD}$, $\mathcal{X} \in \mathfrak{F}$.

Proof. For any filter object \mathcal{Y} we have $\mathcal{Y} \cap^{\mathfrak{F}} \langle f \rangle (\mathcal{X} \cap^{\mathfrak{F}} \text{dom } f) \neq \emptyset \Leftrightarrow \mathcal{X} \cap^{\mathfrak{F}} \text{dom } f \cap^{\mathfrak{F}} \langle f^{-1} \rangle \mathcal{Y} \neq \emptyset \Leftrightarrow \mathcal{X} \cap^{\mathfrak{F}} \text{im } f^{-1} \cap^{\mathfrak{F}} \langle f^{-1} \rangle \mathcal{Y} \neq \emptyset \Leftrightarrow \mathcal{X} \cap^{\mathfrak{F}} \langle f^{-1} \rangle \mathcal{Y} \neq \emptyset \Leftrightarrow \mathcal{Y} \cap^{\mathfrak{F}} \langle f \rangle \mathcal{X} \neq \emptyset$. \square

Proposition 49. $\mathcal{X} \cap^{\mathfrak{F}} \text{dom } f \neq \emptyset \Leftrightarrow \langle f \rangle \mathcal{X} \neq \emptyset$ for any $f \in \text{FCD}$, $\mathcal{X} \in \mathfrak{F}$.

Proof. $\mathcal{X} \cap^{\mathfrak{F}} \text{dom } f \neq \emptyset \Leftrightarrow \mathcal{X} \cap^{\mathfrak{F}} \langle f^{-1} \rangle \mathfrak{U} \neq \emptyset \Leftrightarrow \mathfrak{U} \cap^{\mathfrak{F}} \langle f \rangle \mathcal{X} \neq \emptyset \Leftrightarrow \langle f \rangle \mathcal{X} \neq \emptyset$. \square

Corollary 50. $\text{dom } f = \bigcup^{\mathfrak{F}} \{a \mid a \in \text{atoms}^{\mathfrak{F}} \mathfrak{U}, \langle f \rangle a \neq \emptyset\}$.

Proof. This follows from that \mathfrak{F} is an atomistic lattice. \square

3.7 Category of funcoids

I will define the category FCD of funcoids:

- The set of objects is \mathfrak{F} .
- The set of morphisms from a filter object \mathcal{A} to a filter object \mathcal{B} is the set of triples $(f; \mathcal{A}; \mathcal{B})$ where f is a funcoid such that $\text{dom } f \subseteq \mathcal{A}$, $\text{im } f \subseteq \mathcal{B}$.
- Composition of morphisms is defined in the natural way.
- Identity morphism of a filter object \mathcal{A} is $(I_{\mathcal{A}}; \mathcal{A}; \mathcal{A})$.

To prove that it is really a category is trivial.

3.8 Specifying funcoids by functions or relations on atomic filter objects

Theorem 51. For any funcoid f and filter objects \mathcal{X} and \mathcal{Y}

1. $\langle f \rangle \mathcal{X} = \bigcup^{\mathfrak{F}} \langle \langle f \rangle \rangle \text{atoms}^{\mathfrak{F}} \mathcal{X}$;
2. $\mathcal{X}[f]\mathcal{Y} \Leftrightarrow \exists x \in \text{atoms}^{\mathfrak{F}} \mathcal{X}, y \in \text{atoms}^{\mathfrak{F}} \mathcal{Y}: x[f]y$.

Proof. 1.

$$\begin{aligned} \mathcal{Y} \cap^{\mathfrak{F}} \langle f \rangle \mathcal{X} \neq \emptyset &\Leftrightarrow \mathcal{X} \cap^{\mathfrak{F}} \langle f^{-1} \rangle \mathcal{Y} \neq \emptyset \\ &\Leftrightarrow \exists x \in \text{atoms}^{\mathfrak{F}} \mathcal{X}: x \cap^{\mathfrak{F}} \langle f^{-1} \rangle \mathcal{Y} \neq \emptyset \\ &\Leftrightarrow \exists x \in \text{atoms}^{\mathfrak{F}} \mathcal{X}: \mathcal{Y} \cap^{\mathfrak{F}} \langle f \rangle x \neq \emptyset. \end{aligned}$$

$\partial \langle f \rangle \mathcal{X} = \bigcup \langle \partial \rangle \langle \langle f \rangle \rangle \text{atoms}^{\mathfrak{F}} \mathcal{X} = \partial \bigcup \langle \langle f \rangle \rangle \text{atoms}^{\mathfrak{F}} \mathcal{X}$.

2. If $\mathcal{X}[f]\mathcal{Y}$, then $\mathcal{Y} \cap^{\mathfrak{F}} \langle f \rangle \mathcal{X} \neq \emptyset$, consequently exists $y \in \text{atoms}^{\mathfrak{F}} \mathcal{Y}$ such that $y \cap^{\mathfrak{F}} \langle f \rangle \mathcal{X} \neq \emptyset$, $\mathcal{X}[f]y$. Repeating this second time we get that there exist $x \in \text{atoms}^{\mathfrak{F}} \mathcal{X}$ such that $x[f]y$. From this follows

$$\exists x \in \text{atoms}^{\mathfrak{F}} \mathcal{X}, y \in \text{atoms}^{\mathfrak{F}} \mathcal{Y}: x[f]y.$$

The reverse is obvious. \square

Theorem 52.

1. A function $\alpha \in \mathfrak{F}^{\text{atoms}^{\mathfrak{U}}}$ such that (for any $a \in \text{atoms}^{\mathfrak{U}}$)

$$\alpha a \supseteq \bigcap^{\mathfrak{F}} \left\langle \bigcup^{\mathfrak{F}} \circ \langle \alpha \rangle \circ \text{atoms}^{\mathfrak{F}} \right\rangle \text{up } a \quad (5)$$

can be continued to the function $\langle f \rangle$ for exactly one functor f ;

$$\langle f \rangle \mathcal{X} = \bigcup^{\mathfrak{F}} \langle \alpha \rangle \text{atoms}^{\mathfrak{F}} \mathcal{X} \quad (6)$$

for any filter object \mathcal{X} .

2. A relation $\delta \in \mathcal{P}(\text{atoms}^{\mathfrak{U}})^2$ such that (for any $a, b \in \text{atoms}^{\mathfrak{U}}$)

$$\forall X \in \text{up } a, Y \in \text{up } b \exists x \in \text{atoms}^{\mathfrak{F}} X, y \in \text{atoms}^{\mathfrak{F}} Y: x \delta y \Rightarrow a \delta b \quad (7)$$

can be continued to the relation $[f]$ for exactly one functor f ;

$$\mathcal{X}[f]\mathcal{Y} \Leftrightarrow \exists x \in \text{atoms}^{\mathfrak{F}} \mathcal{X}, y \in \text{atoms}^{\mathfrak{F}} \mathcal{Y}: x \delta y \quad (8)$$

for any filter objects \mathcal{X}, \mathcal{Y} .

Proof. Existence of no more than one such functors and formulas (6) and (8) follow from the previous theorem.

1. Consider the function $\alpha' \in \mathfrak{F}^{\mathfrak{U}}$ defined by the formula (for any $X \in \mathcal{P}\mathfrak{U}$)

$$\alpha' X = \bigcup^{\mathfrak{F}} \langle \alpha \rangle \text{atoms}^{\mathfrak{F}} X.$$

Obviously $\alpha' \emptyset = \emptyset$. For any $I, J \in \mathcal{P}\mathfrak{U}$

$$\begin{aligned} \alpha'(I \cup J) &= \bigcup^{\mathfrak{F}} \langle \alpha' \rangle \text{atoms}^{\mathfrak{F}}(I \cup J) \\ &= \bigcup^{\mathfrak{F}} \langle \alpha' \rangle (\text{atoms}^{\mathfrak{F}} I \cup \text{atoms}^{\mathfrak{F}} J) \\ &= \bigcup^{\mathfrak{F}} (\langle \alpha' \rangle \text{atoms}^{\mathfrak{F}} I \cup \langle \alpha' \rangle \text{atoms}^{\mathfrak{F}} J) \\ &= \bigcup^{\mathfrak{F}} \langle \alpha' \rangle \text{atoms}^{\mathfrak{F}} I \cup \bigcup^{\mathfrak{F}} \langle \alpha' \rangle \text{atoms}^{\mathfrak{F}} J. \\ &= \alpha' I \cup \alpha' J. \end{aligned}$$

Let continue α' till a functor f (by the theorem 25): $\langle f \rangle \mathcal{X} = \bigcap^{\mathfrak{F}} \langle \alpha' \rangle \text{up } \mathcal{X}$.

Let's prove the reverse of (5):

$$\begin{aligned} \bigcap^{\mathfrak{F}} \left\langle \bigcup^{\mathfrak{F}} \circ \langle \alpha \rangle \circ \text{atoms}^{\mathfrak{F}} \right\rangle \text{up } a &= \bigcap^{\mathfrak{F}} \left\langle \bigcup^{\mathfrak{F}} \circ \langle \alpha \rangle \right\rangle \langle \text{atoms}^{\mathfrak{F}} \rangle \text{up } a \\ &\supseteq \bigcap^{\mathfrak{F}} \left\langle \bigcup^{\mathfrak{F}} \circ \langle \alpha \rangle \right\rangle \{ \{ a \} \} \\ &= \bigcap^{\mathfrak{F}} \left\{ \left(\bigcup^{\mathfrak{F}} \circ \langle \alpha \rangle \right) \{ a \} \right\} \\ &= \bigcap^{\mathfrak{F}} \left\{ \bigcup^{\mathfrak{F}} \langle \alpha \rangle \{ a \} \right\} \\ &= \bigcap^{\mathfrak{F}} \left\{ \bigcup^{\mathfrak{F}} \{ \alpha a \} \right\} = \bigcap^{\mathfrak{F}} \{ \alpha a \} = \alpha a. \end{aligned}$$

Finally,

$$\alpha a = \bigcap^{\mathfrak{F}} \left\langle \bigcup^{\mathfrak{F}} \circ \langle \alpha \rangle \circ \text{atoms}^{\mathfrak{F}} \right\rangle \text{up } a = \bigcap^{\mathfrak{F}} \langle \alpha' \rangle \text{up } a = \langle f \rangle a,$$

so $\langle f \rangle$ is a continuation of α .

2. Consider the relation $\delta' \in \mathcal{P}(\mathcal{P}\mathfrak{U})^2$ defined by the formula (for any $X, Y \in \mathcal{P}\mathfrak{U}$)

$$X \delta' Y \Leftrightarrow \exists x \in \text{atoms}^{\mathfrak{F}} X, y \in \text{atoms}^{\mathfrak{F}} Y: x \delta y.$$

Obviously $\neg(X \delta' \emptyset)$ and $\neg(\emptyset \delta' Y)$.

$$\begin{aligned} (I \cup J) \delta' Y &\Leftrightarrow \exists x \in \text{atoms}^{\mathfrak{F}}(I \cup J), y \in \text{atoms}^{\mathfrak{F}} Y: x \delta y \\ &\Leftrightarrow \exists x \in \text{atoms}^{\mathfrak{F}} I \cup \text{atoms}^{\mathfrak{F}} J, y \in \text{atoms}^{\mathfrak{F}} Y: x \delta y \\ &\Leftrightarrow \exists x \in \text{atoms}^{\mathfrak{F}} I, y \in \text{atoms}^{\mathfrak{F}} Y: x \delta y \vee \exists x \in \text{atoms}^{\mathfrak{F}} J, y \in \text{atoms}^{\mathfrak{F}} Y: x \delta y \\ &\Leftrightarrow I \delta' Y \vee J \delta' Y; \end{aligned}$$

analogously $X \delta' (I \cup J) \Leftrightarrow X \delta' I \vee X \delta' J$. Let's continue δ' till a funcoid f (by the theorem 25):

$$\mathcal{X}[f]\mathcal{Y} \Leftrightarrow \forall X \in \text{up } \mathcal{X}, Y \in \text{up } \mathcal{Y}: X \delta' Y$$

The reverse of (7) implication is trivial, so

$$\forall X \in \text{up } a, Y \in \text{up } b \exists x \in \text{atoms}^{\delta} X, y \in \text{atoms}^{\delta} Y: x \delta y \Leftrightarrow a \delta b.$$

$$\forall X \in \text{up } a, Y \in \text{up } b \exists x \in \text{atoms}^{\delta} X, y \in \text{atoms}^{\delta} Y: x \delta y \Leftrightarrow \forall X \in \text{up } a, Y \in \text{up } b: X \delta' Y \Leftrightarrow a[f]b.$$

So $a \delta b \Leftrightarrow a[f]b$, that is $[f]$ is a continuation of δ . \square

One of uses of the previous theorem is proof of the following theorem:

Theorem 53. If R is a set of funcoids, $x, y \in \text{atoms}^{\delta} \mathcal{U}$, then

1. $\langle \bigcap^{\text{FCD}} R \rangle x = \bigcap^{\delta} \{ \langle f \rangle x \mid f \in R \}$;
2. $x[\bigcap^{\text{FCD}} R]y \Leftrightarrow \forall f \in R: x[f]y$.

Proof. 2. Let denote $x \delta y \Leftrightarrow \forall f \in R: x[f]y$.

$$\begin{aligned} \forall X \in \text{up } a, Y \in \text{up } b \exists x \in \text{atoms}^{\delta} X, y \in \text{atoms}^{\delta} Y: x \delta y &\Leftrightarrow \\ \forall f \in R, X \in \text{up } a, Y \in \text{up } b \exists x \in \text{atoms}^{\delta} X, y \in \text{atoms}^{\delta} Y: x[f]y &\Rightarrow \\ \forall f \in R, X \in \text{up } a, Y \in \text{up } b: X[f]Y &\Rightarrow \\ \forall f \in R: a[f]b &\Leftrightarrow \\ a \delta b. & \end{aligned}$$

So, by the theorem 52, δ can be continued till $[p]$ for some funcoid p .

For any funcoid q such that $\forall f \in R: q \subseteq f$ we have $x[q]y \Rightarrow \forall f \in R: x[f]y \Leftrightarrow x \delta y \Leftrightarrow x[p]y$, so $q \subseteq f$. Consequently $p = \bigcap^{\text{FCD}} R$.

From this $x[\bigcap^{\text{FCD}} R]y \Leftrightarrow \forall f \in R: x[f]y$.

1. From the former $y \cap^{\delta} \langle \bigcap^{\text{FCD}} R \rangle x \neq \emptyset \Leftrightarrow \forall f \in R: y \cap^{\delta} \langle f \rangle x \neq \emptyset$ for any $y \in \text{atoms}^{\delta} \mathcal{U}$. From this follows what we need to prove. \square

3.9 Direct product of filter objects

A generalization of direct (Cartesian) product of two sets is direct product of two filter objects as defined in the theory of funcoids:

Definition 54. *Direct product* of filter objects \mathcal{A} and \mathcal{B} is such a funcoid $\mathcal{A} \times^{\text{FCD}} \mathcal{B}$ that

$$\mathcal{X}[\mathcal{A} \times^{\text{FCD}} \mathcal{B}]\mathcal{Y} \Leftrightarrow \mathcal{X} \cap^{\delta} \mathcal{A} \neq \emptyset \wedge \mathcal{Y} \cap^{\delta} \mathcal{B} \neq \emptyset.$$

Proposition 55. $\mathcal{A} \times^{\text{FCD}} \mathcal{B}$ is really a funcoid and

$$\langle \mathcal{A} \times^{\text{FCD}} \mathcal{B} \rangle \mathcal{X} = \begin{cases} \mathcal{B} & \text{if } \mathcal{X} \cap^{\delta} \mathcal{A} \neq \emptyset; \\ \emptyset & \text{if } \mathcal{X} \cap^{\delta} \mathcal{A} = \emptyset. \end{cases}$$

Proof. Obvious. \square

Obvious 56. $A \times B = A \times^{\text{FCD}} B$ for sets A and B .

Proposition 57. $f \subseteq \mathcal{A} \times^{\text{FCD}} \mathcal{B} \Leftrightarrow \text{dom } f \subseteq \mathcal{A} \wedge \text{im } f \subseteq \mathcal{B}$ for any $f \in \text{FCD}$ and $\mathcal{A}, \mathcal{B} \in \mathfrak{F}$.

Proof. If $f \subseteq \mathcal{A} \times^{\text{FCD}} \mathcal{B}$ then $\text{dom } f \subseteq \text{dom}(\mathcal{A} \times^{\text{FCD}} \mathcal{B}) \subseteq \mathcal{A}$, $\text{im } f \subseteq \text{im}(\mathcal{A} \times^{\text{FCD}} \mathcal{B}) \subseteq \mathcal{B}$. If $\text{dom } f \subseteq \mathcal{A} \wedge \text{im } f \subseteq \mathcal{B}$ then

$$\forall \mathcal{X}, \mathcal{Y} \in \mathfrak{F}: (\mathcal{X}[f]\mathcal{Y} \Rightarrow \mathcal{X} \cap^{\delta} \mathcal{A} \neq \emptyset \wedge \mathcal{Y} \cap^{\delta} \mathcal{B} \neq \emptyset);$$

consequently $f \subseteq \mathcal{A} \times^{\text{FCD}} \mathcal{B}$. \square

The following theorem gives a formula for calculating an important particular case of intersection on the lattice of funcoids:

Theorem 58. $f \cap^{\text{FCD}} \mathcal{A} \times^{\text{FCD}} \mathcal{B} = I_B \circ f \circ I_A$ for any $f \in \text{FCD}$ and $\mathcal{A}, \mathcal{B} \in \mathfrak{F}$.

Proof. $h \stackrel{\text{def}}{=} I_B \circ f \circ I_A$. For any $\mathcal{X} \in \mathfrak{F}$

$$\langle h \rangle \mathcal{X} = \langle I_B \rangle \langle f \rangle \langle I_A \rangle \mathcal{X} = \mathcal{B} \cap \langle f \rangle (\mathcal{A} \cap \mathcal{X}).$$

From this, as easy to show, $h \subseteq f$ and $h \subseteq \mathcal{A} \times \mathcal{B}$. If $g \subseteq f \wedge g \subseteq \mathcal{A} \times^{\text{FCD}} \mathcal{B}$ for a funcoid g then $\text{dom } g \subseteq \mathcal{A}$, $\text{im } g \subseteq \mathcal{B}$,

$$\langle g \rangle \mathcal{X} = \mathcal{B} \cap^{\mathfrak{F}} \langle g \rangle (\mathcal{A} \cap^{\mathfrak{F}} \mathcal{X}) \subseteq \mathcal{B} \cap^{\mathfrak{F}} \langle f \rangle (\mathcal{A} \cap^{\mathfrak{F}} \mathcal{X}) = \langle I_B \rangle \langle f \rangle \langle I_A \rangle \mathcal{X} = \langle h \rangle \mathcal{X},$$

$g \subseteq h$. So $h = f \cap^{\text{FCD}} \mathcal{A} \times^{\text{FCD}} \mathcal{B}$. □

Corollary 59. $f|_{\mathcal{A}} = f \cap \mathcal{A} \times^{\text{FCD}} \mathfrak{U}$ for any $f \in \text{FCD}$ and $\mathcal{A} \in \mathfrak{F}$.

Proof. $f \cap^{\text{FCD}} \mathcal{A} \times^{\text{FCD}} \mathfrak{U} = I_{\mathfrak{U}} \circ f \circ I_{\mathcal{A}} = f \circ I_{\mathcal{A}} = f|_{\mathcal{A}}$. □

Corollary 60. $f \cap^{\text{FCD}} (\mathcal{A} \times^{\text{FCD}} \mathcal{B}) \neq \emptyset \Leftrightarrow \mathcal{A}[f]\mathcal{B}$ for any $f \in \text{FCD}$, $\mathcal{A}, \mathcal{B} \in \mathfrak{F}$.

Proof. $f \cap^{\text{FCD}} \mathcal{A} \times^{\text{FCD}} \mathcal{B} \neq \emptyset \Leftrightarrow \langle f \cap^{\text{FCD}} \mathcal{A} \times^{\text{FCD}} \mathcal{B} \rangle \mathfrak{U} \neq \emptyset \Leftrightarrow \langle I_B \circ f \circ I_A \rangle \mathfrak{U} \neq \emptyset \Leftrightarrow \langle I_B \rangle \langle f \rangle \langle I_A \rangle \mathfrak{U} \neq \emptyset \Leftrightarrow \mathcal{B} \cap^{\text{FCD}} \langle f \rangle (\mathcal{A} \cap^{\mathfrak{F}} \mathfrak{U}) \neq \emptyset \Leftrightarrow \mathcal{B} \cap^{\mathfrak{F}} \langle f \rangle \mathcal{A} \neq \emptyset \Leftrightarrow \mathcal{A}[f]\mathcal{B}$. □

Corollary 61. The filtrator of funcoids is star-separable.

Proof. The set of direct products of sets is a separation subset of the lattice of funcoids. □

Theorem 62. If $S \in \mathcal{P}\mathfrak{F}^2$ then

$$\bigcap^{\text{FCD}} \{ \mathcal{A} \times^{\text{FCD}} \mathcal{B} \mid (\mathcal{A}; \mathcal{B}) \in S \} = \bigcap^{\mathfrak{F}} \text{dom } S \times^{\text{FCD}} \bigcap^{\mathfrak{F}} \text{im } S.$$

Proof. If $x \in \text{atoms}^{\mathfrak{F}} \mathfrak{U}$ then by the theorem 53

$$\left\langle \bigcap^{\text{FCD}} \{ \mathcal{A} \times^{\text{FCD}} \mathcal{B} \mid (\mathcal{A}; \mathcal{B}) \in S \} \right\rangle x = \bigcap^{\mathfrak{F}} \{ \langle \mathcal{A} \times^{\text{FCD}} \mathcal{B} \rangle x \mid (\mathcal{A}; \mathcal{B}) \in S \}.$$

If $x \cap^{\mathfrak{F}} \bigcap^{\mathfrak{F}} \text{dom } S \neq \emptyset$ then

$$\begin{aligned} \forall (\mathcal{A}; \mathcal{B}) \in S: (x \cap^{\mathfrak{F}} \mathcal{A} \neq \emptyset \wedge \langle \mathcal{A} \times^{\text{FCD}} \mathcal{B} \rangle x = \mathcal{B}); \\ \{ \langle \mathcal{A} \times^{\text{FCD}} \mathcal{B} \rangle x \mid (\mathcal{A}; \mathcal{B}) \in S \} = \text{im } S; \end{aligned}$$

if $x \cap^{\mathfrak{F}} \bigcap^{\mathfrak{F}} \text{dom } S = \emptyset$ then

$$\begin{aligned} \exists (\mathcal{A}; \mathcal{B}) \in S: (x \cap^{\mathfrak{F}} \mathcal{A} = \emptyset \wedge \langle \mathcal{A} \times^{\text{FCD}} \mathcal{B} \rangle x = \emptyset); \\ \{ \langle \mathcal{A} \times^{\text{FCD}} \mathcal{B} \rangle x \mid (\mathcal{A}; \mathcal{B}) \in S \} \ni \emptyset. \end{aligned}$$

So

$$\left\langle \bigcap^{\text{FCD}} \{ \mathcal{A} \times^{\text{FCD}} \mathcal{B} \mid (\mathcal{A}; \mathcal{B}) \in S \} \right\rangle x = \begin{cases} \bigcap^{\mathfrak{F}} \text{im } S & \text{if } x \cap^{\mathfrak{F}} \bigcap^{\mathfrak{F}} \text{dom } S \neq \emptyset; \\ \emptyset & \text{if } x \cap^{\mathfrak{F}} \bigcap^{\mathfrak{F}} \text{dom } S = \emptyset. \end{cases}$$

From this follows the statement of the theorem. □

Corollary 63. $\mathcal{A}_0 \times^{\text{FCD}} \mathcal{B}_0 \cap^{\text{FCD}} \mathcal{A}_1 \times^{\text{FCD}} \mathcal{B}_1 = (\mathcal{A}_0 \cap^{\text{FCD}} \mathcal{A}_1) \times^{\text{FCD}} (\mathcal{B}_0 \cap^{\text{FCD}} \mathcal{B}_1)$ for any $\mathcal{A}_0, \mathcal{A}_1, \mathcal{B}_0, \mathcal{B}_1 \in \mathfrak{F}$.

Proof. $\mathcal{A}_0 \times^{\text{FCD}} \mathcal{B}_0 \cap^{\text{FCD}} \mathcal{A}_1 \times^{\text{FCD}} \mathcal{B}_1 = \bigcap^{\mathfrak{F}} \{ \mathcal{A}_0 \times^{\text{FCD}} \mathcal{B}_0, \mathcal{A}_1 \times^{\text{FCD}} \mathcal{B}_1 \}$ what is by the last theorem equal to $(\mathcal{A}_0 \cap^{\text{FCD}} \mathcal{A}_1) \times^{\text{FCD}} (\mathcal{B}_0 \cap^{\text{FCD}} \mathcal{B}_1)$. □

Theorem 64. If $\mathcal{A} \in \mathfrak{F}$ then $\mathcal{A} \times^{\text{FCD}}$ is a complete homomorphism of the lattice \mathfrak{F} to a complete sublattice of the lattice FCD, if also $\mathcal{A} \neq \emptyset$ then it is an isomorphism.

Proof. Let $S \in \mathcal{P}\mathfrak{F}$, $X \in \mathcal{P}\mathcal{U}$, $x \in \text{atoms}^{\mathfrak{F}}\mathcal{U}$.

$$\begin{aligned} \left\langle \bigcup^{\text{FCD}} \langle \mathcal{A} \times^{\text{FCD}} \rangle S \right\rangle X &= \bigcup^{\mathfrak{F}} \{ \langle \mathcal{A} \times^{\text{FCD}} \mathcal{B} \rangle X \mid \mathcal{B} \in S \} \\ &= \begin{cases} \bigcup^{\mathfrak{F}} S & \text{if } X \cap^{\mathfrak{F}} \mathcal{A} \neq \emptyset \\ \emptyset & \text{if } X \cap^{\mathfrak{F}} \mathcal{A} = \emptyset \end{cases} \\ &= \langle \mathcal{A} \times^{\text{FCD}} \bigcup^{\mathfrak{F}} S \rangle X; \\ \left\langle \bigcap^{\text{FCD}} \langle \mathcal{A} \times^{\text{FCD}} \rangle S \right\rangle x &= \bigcap^{\mathfrak{F}} \{ \langle \mathcal{A} \times^{\text{FCD}} \mathcal{B} \rangle x \mid \mathcal{B} \in S \} \\ &= \begin{cases} \bigcap^{\mathfrak{F}} S & \text{if } x \cap^{\mathfrak{F}} \mathcal{A} \neq \emptyset \\ \emptyset & \text{if } x \cap^{\mathfrak{F}} \mathcal{A} = \emptyset \end{cases} \\ &= \langle \mathcal{A} \times^{\text{FCD}} \bigcap^{\mathfrak{F}} S \rangle x. \end{aligned}$$

If $\mathcal{A} \neq \emptyset$ then obviously the function $\mathcal{A} \times^{\text{FCD}}$ is injective. \square

The following proposition states that cutting a rectangle of atomic width from a funcoid always produces a rectangular (representable as a direct product of filter objects) funcoid (of atomic width).

Proposition 65. If a is an atomic filter object, $f \in \text{FCD}$ then $f|_a = a \times^{\text{FCD}} \langle f \rangle a$.

Proof. Let $\mathcal{X} \in \mathfrak{F}$.

$$\mathcal{X} \cap^{\mathfrak{F}} a \neq \emptyset \Rightarrow \langle f|_a \rangle \mathcal{X} = \langle f \rangle a, \quad \mathcal{X} \cap^{\mathfrak{F}} a = \emptyset \Rightarrow \langle f|_a \rangle \mathcal{X} = \emptyset. \quad \square$$

3.10 Atomic funcoids

Theorem 66. A funcoid is an atom of the lattice of funcoids iff it is direct product of two atomic filter objects.

Proof.

Direct implication. Let f is an atomic funcoid. Let's get elements $a \in \text{atoms}^{\mathfrak{F}} \text{dom } f$ and $b \in \text{atoms}^{\mathfrak{F}} \langle f \rangle a$. Then for any $\mathcal{X} \in \mathfrak{F}$

$$\mathcal{X} \cap^{\mathfrak{F}} a = \emptyset \Rightarrow \langle a \times^{\text{FCD}} b \rangle \mathcal{X} = \emptyset \subseteq \langle f \rangle \mathcal{X}, \quad \mathcal{X} \cap^{\mathfrak{F}} a \neq \emptyset \Rightarrow \langle a \times^{\text{FCD}} b \rangle \mathcal{X} = b \subseteq \langle f \rangle \mathcal{X}.$$

So $a \times^{\text{FCD}} b \subseteq f$; because f is an atomic funcoid $f = a \times^{\text{FCD}} b$.

Reverse implication. Let $a, b \in \text{atoms}^{\mathfrak{F}}\mathcal{U}$, $f \in \text{FCD}$. If $b \cap^{\mathfrak{F}} \langle f \rangle a = \emptyset$ then $\neg(a[f]b)$, $f \cap^{\text{FCD}} a \times^{\text{FCD}} b = \emptyset$; if $b \subseteq \langle f \rangle a$ then $\forall \mathcal{X} \in \mathfrak{F}: (\mathcal{X} \cap^{\mathfrak{F}} a \neq \emptyset \Rightarrow \langle f \rangle \mathcal{X} \supseteq b)$, $f \supseteq a \times^{\text{FCD}} b$. Consequently $f \cap^{\text{FCD}} a \times^{\text{FCD}} b = \emptyset \vee f \supseteq a \times^{\text{FCD}} b$; that is $a \times^{\text{FCD}} b$ is an atomic filter object. \square

Theorem 67. The lattice of funcoids is atomic.

Proof. Let f is a non-empty funcoid. Then $\text{dom } f \neq \emptyset$, thus by the theorem 46 in [5] exists $a \in \text{atoms } \text{dom } f$. So $\langle f \rangle a \neq \emptyset$ thus exists $b \in \text{atoms } \langle f \rangle a$. Finally the atomic funcoid $a \times^{\text{FCD}} b \subseteq f$. \square

Theorem 68. The lattice of funcoids is atomically separable.

Proof. Let $f, g \in \text{FCD}$, $f \subset g$. Then exists $a \in \text{atoms}^{\mathfrak{F}}\mathcal{U}$ such that $\langle f \rangle a \subset \langle g \rangle a$. So because the lattice \mathfrak{F} is atomically separable then exists $b \in \text{atoms}^{\mathfrak{F}}\mathcal{U}$ such that $\langle f \rangle a \cap^{\mathfrak{F}} b = \emptyset$ and $b \subseteq \langle g \rangle a$. For any $x \in \text{atoms}^{\mathfrak{F}}\mathcal{U}$

$$\begin{aligned} \langle f \rangle a \cap^{\mathfrak{F}} \langle a \times^{\text{FCD}} b \rangle a &= \langle f \rangle a \cap^{\mathfrak{F}} b = \emptyset, \\ x \neq a &\Rightarrow \langle f \rangle x \cap^{\mathfrak{F}} \langle a \times^{\text{FCD}} b \rangle x = \langle f \rangle x \cap^{\mathfrak{F}} \emptyset = \emptyset \end{aligned}$$

Thus $\langle f \rangle x \cap^{\mathfrak{S}} \langle a \times b \rangle x = \emptyset$ and consequently $f \cap^{\text{FCD}} a \times^{\text{FCD}} b = \emptyset$.

$$\begin{aligned} \langle a \times^{\text{FCD}} b \rangle a = b &\subseteq \langle g \rangle a, \\ x \neq a &\Rightarrow \langle a \times^{\text{FCD}} b \rangle x = \emptyset \subseteq \langle g \rangle a. \end{aligned}$$

Thus $\langle a \times^{\text{FCD}} b \rangle x = b \subseteq \langle g \rangle x$ and consequently $a \times^{\text{FCD}} b \subseteq g$.

So the lattice of funcoids is separable by the theorem 19 in [5]. \square

Corollary 69. The lattice of funcoids is:

1. separable;
2. atomically separable;
3. conforming to Wallman's disjunction property.

Proof. By the theorem 22 in [5]. \square

Remark 70. For more ways to characterize (atomic) separability of the lattice of funcoids see [5], subsections "Separation subsets and full stars" and "Atomically separable lattices".

Corollary 71. The lattice of funcoids is an atomistic lattice.

Proof. Let f is a funcoid. Suppose contrary to the statement to be proved that $\bigcup^{\mathfrak{S}} \text{atoms}^{\text{FCD}} f \subset f$. Then exists $a \in \text{atoms}^{\text{FCD}} f$ such that $a \cap^{\mathfrak{S}} \bigcup^{\mathfrak{S}} \text{atoms}^{\text{FCD}} f = \emptyset$ what is impossible. \square

Proposition 72. $\text{atoms}^{\text{FCD}}(f \cup^{\mathfrak{S}} g) = \text{atoms}^{\text{FCD}} f \cup \text{atoms}^{\text{FCD}} g$ for any funcoids f and g .

Proof. $a \times^{\text{FCD}} b \cap^{\text{FCD}} (f \cup^{\text{FCD}} g) \neq \emptyset \Leftrightarrow a[f \cup^{\text{FCD}} g]b \Leftrightarrow a[f]b \vee a[g]b \Leftrightarrow a \times^{\text{FCD}} b \cap^{\text{FCD}} f \neq \emptyset \vee a \times^{\text{FCD}} b \cap^{\text{FCD}} g \neq \emptyset$ for any atomic filter objects a and b . \square

Corollary 73. For any $f, g, h \in \text{FCD}$, $R \in \mathcal{P}\text{FCD}$

1. $f \cap^{\text{FCD}} (g \cup^{\text{FCD}} h) = (f \cap^{\text{FCD}} g) \cup^{\text{FCD}} (f \cap^{\text{FCD}} h)$;
2. $f \cup^{\text{FCD}} \bigcap^{\text{FCD}} R = \bigcap^{\text{FCD}} \langle f \cup^{\text{FCD}} \rangle R$.

Proof. We will take in account that the lattice of funcoids is an atomistic lattice. To be concise I will write atoms instead of $\text{atoms}^{\text{FCD}}$ and \cap and \cup instead of \cap^{FCD} and \cup^{FCD} .

1. $\text{atoms}(f \cap (g \cup h)) = \text{atoms } f \cap \text{atoms}(g \cup h) = \text{atoms } f \cap (\text{atoms } g \cup \text{atoms } h) = (\text{atoms } f \cap \text{atoms } g) \cup (\text{atoms } f \cap \text{atoms } h) = \text{atoms}(f \cap g) \cup \text{atoms}(f \cap h) = \text{atoms}((f \cap g) \cup (f \cap h))$.
2. $\text{atoms}(f \cup \bigcap^{\text{FCD}} R) = \text{atoms } f \cup \text{atoms} \bigcap^{\text{FCD}} R = \text{atoms } f \cup \bigcap^{\text{FCD}} \langle \text{atoms} \rangle R = \bigcap^{\text{FCD}} \langle (\text{atoms } f) \cup \rangle \langle \text{atoms} \rangle R = \bigcap^{\text{FCD}} \langle \text{atoms} \rangle \langle f \cup \rangle R = \text{atoms} \bigcap^{\text{FCD}} \langle f \cup \rangle R$. \square

Note that distributivity of the lattice of funcoids is proved through using atoms of this lattice. I have never seen such method of proving distributivity.

The next proposition is one more (among the theorem 42) generalization for funcoids of composition of relations.

Corollary 74. The lattice of funcoids is co-brouwerian.

Proposition 75. For any $f, g \in \text{FCD}$

$\text{atoms}^{\text{FCD}}(g \circ f) = \{x \times^{\text{FCD}} z \mid x, z \in \text{atoms}^{\mathfrak{S}} \mathcal{U}, \exists y \in \text{atoms}^{\mathfrak{S}} \mathcal{U}: (x \times^{\text{FCD}} y \in \text{atoms}^{\text{FCD}} f \wedge y \times^{\text{FCD}} z \in \text{atoms}^{\text{FCD}} g)\}$.

Proof. $x \times^{\text{FCD}} z \cap^{\text{FCD}} g \circ f \neq \emptyset \Leftrightarrow x[g \circ f]z \Leftrightarrow \exists y \in \text{atoms}^{\mathfrak{S}} \mathcal{U}: (x[f]y \wedge y[g]z) \Leftrightarrow \exists y \in \text{atoms}^{\mathfrak{S}} \mathcal{U}: (x \times^{\text{FCD}} y \cap^{\text{FCD}} f \neq \emptyset \wedge y \times^{\text{FCD}} z \cap^{\text{FCD}} g \neq \emptyset)$ (were used the theorem 42). \square

Conjecture 76. The set of discrete funcoids is the center of the lattice of funcoids.

3.11 Complete funcoids

Definition 77. I will call *co-complete* such a funcoid f that $\forall X \in \mathcal{P}\mathcal{U}: \langle f \rangle X \in \mathcal{P}\mathcal{U}$.

Remark 78. I will call *generalized closure* such a function $\alpha \in \mathcal{P}\mathcal{U}^{\mathcal{P}\mathcal{U}}$ that

1. $\alpha \emptyset = \emptyset$;
2. $\forall I, J \in \mathcal{P}\mathcal{U}: \alpha(I \cup J) = \alpha I \cup^{\mathfrak{F}} \alpha J$.

Obvious 79. A funcoid f is co-complete iff $\langle f \rangle|_{\mathcal{P}\mathcal{U}}$ is a generalized closure.

Remark 80. Thus funcoids can be considered as a generalization of generalized closures. A topological space in Kuratowski sense is the same as reflexive and transitive generalized closure. So topological spaces can be considered as a special case of funcoids.

Definition 81. I will call a *complete funcoid* a funcoid whose reverse is co-complete.

Theorem 82. The following conditions are equivalent for every funcoid f :

1. funcoid f is complete.
2. $\forall S \in \mathcal{P}\mathfrak{F}, J \in \mathcal{P}\mathcal{U}: (\bigcup^{\mathfrak{F}} S[f]J \Leftrightarrow \exists I \in S: I[f]J)$;
3. $\forall S \in \mathcal{P}\mathcal{P}\mathcal{U}, J \in \mathcal{P}\mathcal{U}: (\bigcup S[f]J \Leftrightarrow \exists I \in S: I[f]J)$;
4. $\forall S \in \mathcal{P}\mathfrak{F}: \langle f \rangle \bigcup^{\mathfrak{F}} S = \bigcup^{\mathfrak{F}} \langle \langle f \rangle \rangle S$;
5. $\forall S \in \mathcal{P}\mathcal{P}\mathcal{U}: \langle f \rangle \bigcup S = \bigcup^{\mathfrak{F}} \langle \langle f \rangle \rangle S$;
6. $\forall A \in \mathcal{P}\mathcal{U}: \langle f \rangle A = \bigcup^{\mathfrak{F}} \{ \langle f \rangle a \mid a \in A \}$.

Proof.

(3) \Rightarrow (1). For any $S \in \mathcal{P}\mathcal{P}\mathcal{U}, J \in \mathcal{P}\mathcal{U}$

$$\bigcup S \cap^{\mathfrak{F}} \langle f^{-1} \rangle J \neq \emptyset \Leftrightarrow \exists I \in S: I \cap^{\mathfrak{F}} \langle f^{-1} \rangle J \neq \emptyset, \quad (9)$$

consequently by the theorem 52 in [5] we have $\langle f^{-1} \rangle J \in \mathcal{P}\mathcal{U}$.

(1) \Rightarrow (2). For any $S \in \mathcal{P}\mathfrak{F}, J \in \mathcal{P}\mathcal{U}$ we have $\langle f^{-1} \rangle J \in \mathcal{P}\mathcal{U}$, consequently the formula (9) is true. From this follows (2).

(6) \Rightarrow (5). $\langle f \rangle \bigcup S = \bigcup^{\mathfrak{F}} \{ \langle f \rangle a \mid a \in \bigcup S \} = \bigcup^{\mathfrak{F}} \{ \bigcup^{\mathfrak{F}} \{ \langle f \rangle a \mid a \in A \} \mid A \in S \} = \bigcup^{\mathfrak{F}} \{ \langle f \rangle A \mid A \in S \} = \bigcup^{\mathfrak{F}} \langle \langle f \rangle \rangle S$.

(2) \Rightarrow (3), (4) \Rightarrow (5), (5) \Rightarrow (3), (2) \Rightarrow (4), (5) \Rightarrow (6). Obvious. \square

The following proposition shows that complete funcoids are a direct generalization of pre-topological spaces.

Proposition 83. To specify a complete funcoid f it is enough to specify $\langle f \rangle$ on one-element sets, values of $\langle f \rangle$ on one element sets can be specified arbitrarily.

Proof. From the above theorem is clear that knowing $\langle f \rangle$ on one-element sets $\langle f \rangle$ can be found on any sets and then its value can be inferred for any filter objects.

Choosing arbitrarily the values of $\langle f \rangle$ on one-element sets we can define a complete funcoid the following way: $\langle f \rangle X \stackrel{\text{def}}{=} \bigcup^{\mathfrak{F}} \{ \langle f \rangle \{ \alpha \} \mid \alpha \in X \}$ for any $X \in \mathcal{P}\mathcal{U}$. Obviously it is really a complete funcoid. \square

Theorem 84. A funcoid is discrete iff it is both complete and co-complete.

Proof.

Direct implication. Obvious.

Reverse implication. Let f is both a complete and co-complete funcoïd. Consider the relation g defined by that $\langle g \rangle \{ \alpha \} = \langle f \rangle \{ \alpha \}$ (g is correctly defined because f is a generalized closure). Because f is a complete funcoïd $f = g$. \square

Theorem 85. If R is a set of (co-)complete funcoïds then $\bigcup^{\text{FCD}} R$ is a (co-)complete funcoïd.

Proof. It is enough to prove only for co-complete funcoïds. Let R is a set of co-complete funcoïds. Then for any $X \in \mathcal{P}\mathcal{U}$

$$\left\langle \bigcup^{\text{FCD}} R \right\rangle X = \bigcup \{ \langle f \rangle X \mid f \in R \} \in \mathcal{P}\mathcal{U}$$

(used the theorem 39). \square

Corollary 86. If R is a set of binary relations then $\bigcup^{\text{FCD}} R = \bigcup R$.

Proof. From two last theorems. \square

Theorem 87. The filtrator of funcoïds is filtered.

Proof. It's enough to prove that every funcoïd is representable as (infinite) intersection (on the lattice of funcoïds) of some set of discrete funcoïds.

Let $f \in \text{FCD}$, $A \in \mathcal{P}\mathcal{U}$, $B \in \text{up}\langle f \rangle A$, $g(A; B) \stackrel{\text{def}}{=} A \times B \cup^{\text{FCD}} \bar{A} \times \mathcal{U}$. For any $X \in \mathcal{P}\mathcal{U}$

$$\langle g(A; B) \rangle X = \langle A \times^{\text{FCD}} B \rangle X \cup \langle \bar{A} \times^{\text{FCD}} \mathcal{U} \rangle X = \left(\begin{array}{l} \emptyset \text{ if } X = \emptyset \\ B \text{ if } \emptyset \neq X \subseteq A \\ \mathcal{U} \text{ if } X \not\subseteq A \end{array} \right) \supseteq \langle f \rangle X;$$

so $g(A; B) \supseteq f$. For any $A \in \mathcal{P}\mathcal{U}$

$$\bigcap^{\mathfrak{S}} \{ \langle g(A; B) \rangle A \mid B \in \text{up}\langle f \rangle A \} = \bigcap^{\mathfrak{S}} \{ B \mid B \in \text{up}\langle f \rangle A \} = \langle f \rangle A;$$

consequently

$$\bigcap^{\text{FCD}} \{ g(A; B) \mid A \in \mathcal{P}\mathcal{U}, B \in \text{up}\langle f \rangle A \} = f. \quad \square$$

In certain cases the theorem 43 can be generalized for infinite unions.

Theorem 88. Let $f \in \text{FCD}$. If R is a set of co-complete funcoïds then

$$f \circ \bigcup^{\text{FCD}} R = \bigcup^{\text{FCD}} \langle f \circ \rangle R.$$

Proof. If R is a set of co-complete funcoïds then for $X, Z \in \mathcal{P}\mathcal{U}$

$$\begin{aligned} X \left[f \circ \bigcup^{\text{FCD}} R \right] Z &\Leftrightarrow \\ \text{(by the theorem 42)} &\Leftrightarrow \exists y \in \text{atoms}^{\mathfrak{S}\mathcal{U}} : \left(X \left[\bigcup^{\text{FCD}} R \right] y \wedge y[f]Z \right) \\ &\Leftrightarrow \exists y \in \text{atoms}^{\mathfrak{S}\mathcal{U}} : \left(y \cap^{\mathfrak{S}} \left\langle \bigcup^{\text{FCD}} R \right\rangle X \neq \emptyset \wedge y[f]Z \right) \\ \text{(by the theorem 39)} &\Leftrightarrow \exists y \in \text{atoms}^{\mathfrak{S}\mathcal{U}} : \left(y \cap^{\mathfrak{S}} \bigcup \{ \langle u \rangle X \mid u \in R \} \neq \emptyset \wedge y[f]Z \right) \\ \text{(by the theorem 52 in [5])} &\Leftrightarrow \exists y \in \text{atoms}^{\mathfrak{S}\mathcal{U}} : \left(\exists u \in R : y \cap^{\mathfrak{S}} \langle u \rangle X \neq \emptyset \wedge y[f]Z \right) \\ &\Leftrightarrow \exists y \in \text{atoms}^{\mathfrak{S}\mathcal{U}} : \left(\exists u \in R : X[u]y \wedge y[f]Z \right) \\ &\Leftrightarrow \exists y \in \text{atoms}^{\mathfrak{S}\mathcal{U}}, u \in R : \left(X[u]y \wedge y[f]Z \right) \\ \text{(by the theorem 42)} &\Leftrightarrow \exists u \in R : X[f \circ u]Z \\ \text{(by the theorem 39)} &\Leftrightarrow X \left[\bigcup^{\text{FCD}} \langle f \circ \rangle R \right] Z. \end{aligned}$$

\square

3.12 Completion of functors

I will denote ComplFCD and CoComplFCD the sets of complete and co-complete functors correspondingly.

Obvious 89. ComplFCD and CoComplFCD are closed regarding composition of functors.

Proposition 90. ComplFCD and CoComplFCD (with induced order) are complete lattices.

Proof. Follows from the corollary 85. □

Theorem 91. $\text{Cor } f = \text{Cor}' f$ for an element f of the filtrator of functors.

Proof. From the theorem 26 in [5] and the corollary 86 and theorem 87. □

Definition 92. *Completion* of a functor f is the complete functor $\text{Compl } f$ defined by the formula $\langle \text{Compl } f \rangle \{ \alpha \} = \langle f \rangle \{ \alpha \}$ for $\alpha \in \mathcal{U}$.

Definition 93. *Co-completion* of a functor f is defined by the formula

$$\text{CoCompl } f = (\text{Compl } f^{-1})^{-1}.$$

Proposition 94. The filtrator $(\text{FCD}; \text{ComplFCD})$ is filtered.

Proof. Because the filtrator $(\text{FCD}; \mathcal{P}\mathcal{U}^2)$ is filtered. □

Theorem 95. $\text{Compl } f = \text{Cor}^{(\text{FCD}; \text{ComplFCD})} f = \text{Cor}'^{(\text{FCD}; \text{ComplFCD})} f$.

Proof. $\text{Cor}^{(\text{FCD}; \text{ComplFCD})} f = \text{Cor}'^{(\text{FCD}; \text{ComplFCD})} f$ since (the theorem 26 in [5]) the filtrator $(\text{FCD}; \text{ComplFCD})$ is filtered (as a consequence of the theorem 87) and with join closed core (the theorem 85).

Let $g \in \text{up}^{(\text{FCD}; \text{ComplFCD})} f$. Then $g \in \text{ComplFCD}$ and $g \supseteq f$. Thus $g = \text{Compl } g \supseteq \text{Compl } f$.

Thus $\forall g \in \text{up}^{(\text{FCD}; \text{ComplFCD})} f: g \supseteq \text{Compl } f$.

Let $\forall g \in \text{up}^{(\text{FCD}; \text{ComplFCD})} f: h \subseteq g$ for some $h \in \text{ComplFCD}$.

Then $h \subseteq \bigcap^{\text{FCD}} \text{up}^{(\text{FCD}; \text{ComplFCD})} f = f$ and consequently $h = \text{Compl } h \subseteq \text{Compl } f$.

Thus $\text{Compl } f = \bigcap^{\text{ComplFCD}} \text{up}^{(\text{FCD}; \text{ComplFCD})} f = \text{Cor}^{(\text{FCD}; \text{ComplFCD})} f$. □

Theorem 96. Atoms of the lattice ComplFCD are exactly direct products of the form $\{ \alpha \} \times^{\text{FCD}} b$ where $\alpha \in \mathcal{U}$ and b is an atomic f.o.

Proof. First, easy to see that $\{ \alpha \} \times^{\text{FCD}} b$ are elements of ComplFCD . Also \emptyset is an element of ComplFCD .

$\{ \alpha \} \times^{\text{FCD}} b$ are atoms of ComplFCD because these are atoms of FCD .

Remain to prove that if f is an atom of ComplFCD then $f = \{ \alpha \} \times^{\text{FCD}} b$ for some $\alpha \in \mathcal{U}$ and an atomic f.o. b .

Suppose f is a non-empty complete functor. Then exists $\alpha \in \mathcal{U}$ such that $\langle f \rangle \alpha \neq \emptyset$. Thus $\{ \alpha \} \times^{\text{FCD}} b \subseteq f$ for some atomic f.o. b . If f is an atom then $f = \{ \alpha \} \times^{\text{FCD}} b$. □

Theorem 97. $\langle \text{CoCompl } f \rangle X = \text{Cor } \langle f \rangle X$ for every functor f and set X .

Proof. $\text{CoCompl } f \subseteq f$ thus $\langle \text{CoCompl } f \rangle X \subseteq \langle f \rangle X$, but $\langle \text{CoCompl } f \rangle X \in \mathcal{P}\mathcal{U}$ thus $\langle \text{CoCompl } f \rangle X \subseteq \text{Cor } \langle f \rangle X$.

Let $\alpha X = \text{Cor } \langle f \rangle X$. Then $h\emptyset = \emptyset$ and

$$\alpha(X \cup Y) = \text{Cor } \langle f \rangle (X \cup Y) = \text{Cor}(\langle f \rangle X \cup \langle f \rangle Y) = \text{Cor } \langle f \rangle X \cup \text{Cor } \langle f \rangle Y = \alpha X \cup \alpha Y.$$

(used the theorem 64 from [5]). Thus α can be continued till $\langle g \rangle$ for some functor g . This functor is co-complete.

Evidently g is the greatest co-complete functor which is lower than f .

Thus $g = \text{CoCompl } f$ and so $\text{Cor } \langle f \rangle X = \alpha X = \langle g \rangle X = \langle \text{CoCompl } f \rangle X$. \square

Theorem 98. $\text{Compl } f \cap^{\text{FCD}} \text{Compl } g = \text{Compl}(f \cap^{\text{FCD}} g)$ for every functors f and g .

Proof. $\langle \text{CoCompl } f \cap^{\text{FCD}} \text{CoCompl } g \rangle x = \langle \text{CoCompl } f \rangle x \cap^{\mathfrak{S}} \langle \text{CoCompl } g \rangle x = \text{Cor } \langle f \rangle x \cap^{\mathfrak{S}} \text{Cor } \langle g \rangle x = \text{Cor } \langle f \rangle x \cap \text{Cor } \langle g \rangle x = \text{Cor}(\langle f \rangle x \cap^{\mathfrak{S}} \langle g \rangle x) = \text{Cor } \langle f \cap^{\text{FCD}} g \rangle x = \langle \text{CoCompl}(f \cap^{\text{FCD}} g) \rangle x$ for every atomic f.o. x (used the theorem 63 from [5]).

Thus $\text{CoCompl } f \cap^{\text{FCD}} \text{CoCompl } g = \text{CoCompl}(f \cap^{\text{FCD}} g)$. \square

Corollary 99. If f, g are (co-)complete functors then $f \cap^{\text{FCD}} g$ is a (co-)complete functor.

Proof. Let f and g are complete functors.

$f \cap^{\text{FCD}} g = \text{Compl } f \cap^{\text{FCD}} \text{Compl } g = \text{Compl}(f \cap^{\text{FCD}} g) \in \text{ComplFCD}$. \square

Corollary 100. If f, g are discrete functors then $f \cap^{\text{FCD}} g$ is a discrete functor.

Proof. Let f, g are discrete functors. Then $f \cap^{\text{FCD}} g$ is both complete and co-complete. \square

Corollary 101. If f, g are discrete functors then $f \cap^{\text{FCD}} g = f \cap g$.

Theorem 102. ComplFCD is an atomistic lattice.

Proof. Let $f \in \text{ComplFCD}$. $\langle f \rangle X = \bigcup^{\mathfrak{S}} \{ \langle f \rangle x \mid x \in X \} = \bigcup^{\mathfrak{S}} \{ \langle f|_{\{x\}} \rangle x \mid x \in X \}$, thus $f = \bigcup^{\text{FCD}} \{ f|_{\{x\}} \mid x \in X \}$. It is trivial that every $f|_{\{x\}}$ is a union of atoms of ComplFCD . \square

Theorem 103. A functor is complete iff it is a join (on the lattice FCD) of atomic complete functors.

Proof. Follows from the theorem 85 and the previous theorem. \square

[TODO: Research the filtrators related to completion: Is ComplFCD a distributive lattice? Is ComplFCD a co-brouwerian lattice? Is it a separable lattice? Is the filtrator ($\text{FCD}; \text{ComplFCD}$) with (finitely) meet/join closed core? Is it with separeble core? Is it star-separable?]

[TODO: Complete reloids can be defined as join of $\{\alpha\} \times b$ atoms. Composition of complete reloids is complete. Complete reloids are convex. Relationships of complete reloids and complete functors.]

Lemma 104. Co-completion of a complete functor is complete.

Proof. Let f is a complete functor.

$\langle \text{CoCompl } f \rangle X = \text{Cor } \langle f \rangle X = \text{Cor } \bigcup^{\mathfrak{S}} \{ f x \mid x \in X \} = \bigcup \{ \text{Cor } f x \mid x \in X \} = \bigcup \{ \langle \text{CoCompl } f \rangle x \mid x \in X \}$. \square

Theorem 105. $\text{Compl } \text{CoCompl } f = \text{CoCompl } \text{Compl } f = \text{Cor } f$ for every functor f .

Proof. $\text{Compl } \text{CoCompl } f$ is co-complete since (used the lemma) $\text{CoCompl } f$ is co-complete. Thus $\text{Compl } \text{CoCompl } f$ is a discrete functor. Trivially [TODO: Detailed proof] it is the greatest discrete functor under f . Thus $\text{Compl } \text{CoCompl } f = \text{Cor } f$. Similarly $\text{CoCompl } \text{Compl } f = \text{Cor } f$. \square

Theorem 106. $\text{Cor } f \cap^{\text{FCD}} \text{Cor } g = \text{Cor}(f \cap^{\text{FCD}} g)$.

Proof. From the previous theorem and the theorem 98. \square

3.13 Monovalued functors

Following the idea of definition of monovalued morphism let's call *monovalued* such a functor f that $f \circ f^{-1} \subseteq I_{\text{im } f}$.

Obvious 107. A morphism $(f; \mathcal{A}; \mathcal{B})$ of the category of functors is monovalued iff the functor f is monovalued.

Theorem 108. The following statements are equivalent for a functor f :

1. f is monovalued.
2. $\forall a \in \text{atoms}^{\mathfrak{S}} \mathcal{A}: \langle f \rangle a \in \text{atoms}^{\mathfrak{S}} \mathcal{U} \cup \{\emptyset\}$.
3. $\forall \mathcal{I}, \mathcal{J} \in \mathfrak{F}: \langle f^{-1} \rangle (\mathcal{I} \cap^{\mathfrak{S}} \mathcal{J}) = \langle f^{-1} \rangle \mathcal{I} \cap^{\mathfrak{S}} \langle f^{-1} \rangle \mathcal{J}$.
4. $\forall I, J \in \mathcal{P}\mathcal{U}: \langle f^{-1} \rangle (I \cap J) = \langle f^{-1} \rangle I \cap^{\mathfrak{S}} \langle f^{-1} \rangle J$.

Proof.

(2) \Rightarrow (3). Let $a \in \text{atoms}^{\mathfrak{S}} \mathcal{U}$, $\langle f \rangle a = b$. Then because $b \in \text{atoms}^{\mathfrak{S}} \mathcal{U} \cup \{\emptyset\}$

$$\begin{aligned} (\mathcal{I} \cap^{\mathfrak{S}} \mathcal{J}) \cap^{\mathfrak{S}} b \neq \emptyset &\Leftrightarrow \mathcal{I} \cap^{\mathfrak{S}} b \neq \emptyset \wedge \mathcal{J} \cap^{\mathfrak{S}} b \neq \emptyset; \\ a[f](\mathcal{I} \cap^{\mathfrak{S}} \mathcal{J}) &\Leftrightarrow a[f]\mathcal{I} \wedge a[f]\mathcal{J}; \\ (\mathcal{I} \cap^{\mathfrak{S}} \mathcal{J})[f^{-1}]a &\Leftrightarrow \mathcal{I}[f^{-1}]a \wedge \mathcal{J}[f^{-1}]a; \\ a \cap^{\mathfrak{S}} \langle f^{-1} \rangle (\mathcal{I} \cap^{\mathfrak{S}} \mathcal{J}) \neq \emptyset &\Leftrightarrow a \cap^{\mathfrak{S}} \langle f^{-1} \rangle \mathcal{I} \neq \emptyset \wedge a \cap^{\mathfrak{S}} \langle f^{-1} \rangle \mathcal{J} \neq \emptyset; \\ \langle f^{-1} \rangle (\mathcal{I} \cap^{\mathfrak{S}} \mathcal{J}) &= \langle f^{-1} \rangle \mathcal{I} \cap^{\mathfrak{S}} \langle f^{-1} \rangle \mathcal{J}. \end{aligned}$$

(4) \Rightarrow (1). $\langle f^{-1} \rangle a \cap^{\mathfrak{S}} \langle f^{-1} \rangle b = \emptyset$ for any two distinct atomic filter objects a and b . This is equivalent to $\neg(b[f^{-1}]\langle f^{-1} \rangle a)$; $\neg(\langle f^{-1} \rangle a[f]b)$; $b \cap^{\mathfrak{S}} \langle f \rangle \langle f^{-1} \rangle a = \emptyset$; $b \cap^{\mathfrak{S}} \langle f \circ f^{-1} \rangle a = \emptyset$; $\neg(a[f \circ f^{-1}]b)$. So $a[f \circ f^{-1}]b \Rightarrow a = b$ for any atomic filter objects a and b . This is possible only when $f \circ f^{-1} \subseteq I_{\text{Dst } f}$.

(3) \Rightarrow (4). Obvious.

\neg (2) \Rightarrow \neg (1). Suppose $\langle f \rangle a \notin \text{atoms}^{\mathfrak{S}} \mathcal{B} \cup \{\emptyset\}$ for some $a \in \text{atoms}^{\mathfrak{S}} \mathcal{A}$. Then there exist two atomic filter objects $p \neq q$ such that $\langle f \rangle a \supseteq p \wedge \langle f \rangle a \supseteq q$. Consequently $p \cap^{\mathfrak{S}} \langle f \rangle a \neq \emptyset$; $a \cap^{\mathfrak{S}} \langle f^{-1} \rangle p \neq \emptyset$; $a \subseteq \langle f^{-1} \rangle p$; $\langle f \circ f^{-1} \rangle p = \langle f \rangle \langle f^{-1} \rangle p \supseteq \langle f \rangle a \supseteq q$; $\langle f \circ f^{-1} \rangle p \not\subseteq p$. So it cannot be $f \circ f^{-1} \subseteq I_{\text{Dst } f}$. \square

Corollary 109. A function is a monovalued functor.

Remark 110. This corollary can be reformulated as follows: For binary relations the classic concept of monovaluedness and monovaluedness in the above defined sense of monovaluedness of a functor are the same.

Proof. Because $\forall I, J \in \mathcal{P}\mathcal{U}: \langle f^{-1} \rangle (I \cap J) = \langle f^{-1} \rangle I \cap^{\mathfrak{S}} \langle f^{-1} \rangle J$ is true for any function f . \square

3.14 T_1 - and T_2 -separable functors

For functors can be generalized T_0 -, T_1 - and T_2 - separability. Worthwhile note that T_0 and T_2 separability is defined through T_1 separability.

Definition 111. Let call T_1 -separable such functor f that for any $\alpha, \beta \in \mathcal{U}$ is true

$$\alpha \neq \beta \Rightarrow \neg(\{\alpha\}[f]\{\beta\})$$

Definition 112. Let call T_0 -separable such functor f that $f \cap^{\text{FCD}} f^{-1}$ is T_1 -separable.

Definition 113. Let call T_2 -separable such functor f that the functor $f^{-1} \circ f$ is T_1 -separable.

For symmetric transitive functors T_1 - and T_2 -separability are the same (see theorem 13).

3.15 Filter objects closed regarding a functor

Definition 114. Let's call *closed* regarding a functor f such filter object \mathcal{A} that $\langle f \rangle \mathcal{A} \subseteq \mathcal{A}$.

This is a generalization of closedness of a set regarding an unary operation.

Proposition 115. If \mathcal{I} and \mathcal{J} are closed (regarding some funcoïd), S is a set of closed filter objects, then

1. $\mathcal{I} \cup^{\mathfrak{F}} \mathcal{J}$ is a closed filter object;
2. $\bigcap^{\mathfrak{F}} S$ is a closed filter object.

Proof. Let denote the given funcoïd as f . $\langle f \rangle(\mathcal{I} \cup^{\mathfrak{F}} \mathcal{J}) = \langle f \rangle \mathcal{I} \cup^{\mathfrak{F}} \langle f \rangle \mathcal{J} \subseteq \mathcal{I} \cup^{\mathfrak{F}} \mathcal{J}$, $\langle f \rangle \bigcap^{\mathfrak{F}} S \subseteq \bigcap^{\mathfrak{F}} \langle f \rangle S \subseteq \bigcap^{\mathfrak{F}} S$. Consequently the filter objects $\mathcal{I} \cup^{\mathfrak{F}} \mathcal{J}$ and $\bigcap^{\mathfrak{F}} S$ are closed. \square

Proposition 116. If S is a set of closed regarding a complete funcoïd filter objects, then the filter object $\bigcup^{\mathfrak{F}} S$ is also closed regarding our funcoïd.

Proof. $\langle f \rangle \bigcup^{\mathfrak{F}} S = \bigcap^{\mathfrak{F}} \langle f \rangle S \subseteq \bigcup^{\mathfrak{F}} S$ where f is the given funcoïd. \square

4 Reloids

Definition 117. I will call a *reloid* a filter object on the set of binary relations.

Reloids are a generalization of uniform spaces. Also reloids are generalization of binary relations.

Definition 118. The *reverse* reloid of a reloid f is defined by the formula

$$\text{up } f^{-1} = \{F^{-1} \mid F \in \text{up } f^{-1}\}.$$

Reverse reloid is a generalization of conjugate quasi-uniformity.

4.1 Composition of reloids

Definition 119. Composition of reloids is defined by the formula

$$g \circ f = \bigcap^{\text{RLD}} \{G \circ F \mid F \in \text{up } f, G \in \text{up } g\}.$$

Composition of reloids is a reloid.

Lemma 120. $(h \circ g) \circ f = h \circ (g \circ f)$ for any reloids f, g, h .

Proof. For two nonempty collections A and B of sets I will denote

$$A \sim B \Leftrightarrow (\forall K \in A \exists L \in B: L \subseteq K) \wedge (\forall K \in B \exists L \in A: L \subseteq K).$$

It is easy to see that \sim is a transitive relation.

I will denote $B \circ A = \{L \circ K \mid K \in A, L \in B\}$.

Let first prove that for any nonempty collections of relations A, B, C

$$A \sim B \Rightarrow A \circ C \sim B \circ C.$$

Suppose $A \sim B$ and $P \in A \circ C$ that is $K \in A$ and $M \in C$ such that $P = K \circ M$. $\exists K' \in B: K' \subseteq K$ because $A \sim B$. We have $P' = K' \circ M \in B \circ C$. Obviously $P' \subseteq P$. So for any $P \in A \circ C$ exist $P' \in B \circ C$ such that $P' \subseteq P$; vice verse is analogous. So $A \circ C \sim B \circ C$.

$\text{up}((h \circ g) \circ f) \sim \text{up}(h \circ g) \circ \text{up } f$, $\text{up}(h \circ g) \sim (\text{up } h) \circ (\text{up } g)$. By proven above $\text{up}((h \circ g) \circ f) \sim (\text{up } h) \circ (\text{up } g) \circ (\text{up } f)$.

Analogously $\text{up}(h \circ (g \circ f)) \sim (\text{up } h) \circ (\text{up } g) \circ (\text{up } f)$.

So $\text{up}((h \circ g) \circ f) \sim \text{up}(h \circ (g \circ f))$ what is possible only if $\text{up}((h \circ g) \circ f) = \text{up}(h \circ (g \circ f))$. \square

Theorem 121.

1. $f \circ f = \bigcap^{\text{RLD}} \{F \circ F \mid F \in \text{up } f\}$;
2. $f^{-1} \circ f = \bigcap^{\text{RLD}} \{F^{-1} \circ F \mid F \in \text{up } f\}$;
3. $f \circ f^{-1} = \bigcap^{\text{RLD}} \{F \circ F^{-1} \mid F \in \text{up } f\}$.

Proof. I will prove only (1) and (2) because (3) is analogous to (2). □

[FIXME: Proof here.]

Conjecture 122. If f, g, h are reloids then

1. $f \circ (g \cup^{\text{RLD}} h) = f \circ g \cup^{\text{RLD}} f \circ h$;
2. $(g \cup^{\text{RLD}} h) \circ f = g \circ f \cup^{\text{RLD}} h \circ f$.

4.2 Direct product of filter objects

In theory of reloids direct product of filter objects \mathcal{A} and \mathcal{B} is defined by the formula

$$\mathcal{A} \times^{\text{RLD}} \mathcal{B} \stackrel{\text{def}}{=} \bigcap^{\tilde{\mathfrak{F}}} \{A \times B \mid A \in \text{up } \mathcal{A}, B \in \text{up } \mathcal{B}\}.$$

Theorem 123. $\mathcal{A} \times^{\text{RLD}} \mathcal{B} = \bigcup^{\tilde{\mathfrak{F}}} \{a \times^{\text{RLD}} b \mid a \in \text{atoms}^{\tilde{\mathfrak{F}}} \mathcal{A}, b \in \text{atoms}^{\tilde{\mathfrak{F}}} \mathcal{B}\}$ for any $\mathcal{A}, \mathcal{B} \in \tilde{\mathfrak{F}}$.

Proof. Obviously

$$\mathcal{A} \times^{\text{RLD}} \mathcal{B} \supseteq \bigcup^{\tilde{\mathfrak{F}}} \{a \times^{\text{RLD}} b \mid a \in \text{atoms}^{\tilde{\mathfrak{F}}} \mathcal{A}, b \in \text{atoms}^{\tilde{\mathfrak{F}}} \mathcal{B}\}$$

Reversely, let $K \in \text{up} \bigcup^{\tilde{\mathfrak{F}}} \{a \times^{\text{RLD}} b \mid a \in \text{atoms}^{\tilde{\mathfrak{F}}} \mathcal{A}, b \in \text{atoms}^{\tilde{\mathfrak{F}}} \mathcal{B}\}$. Then $K \in \text{up}(a \times^{\text{RLD}} b)$ for every $a \in \text{atoms}^{\tilde{\mathfrak{F}}} \mathcal{A}, b \in \text{atoms}^{\tilde{\mathfrak{F}}} \mathcal{B}$; $K \supseteq X_a \times^{\text{RLD}} Y_b$ for some $X_a \in \text{up } a, Y_b \in \text{up } b$; $K \supseteq \bigcup \{X_a \times Y_b \mid a \in \text{atoms}^{\tilde{\mathfrak{F}}} \mathcal{A}, b \in \text{atoms}^{\tilde{\mathfrak{F}}} \mathcal{B}\} = \bigcup \{X_a \mid a \in \text{atoms}^{\tilde{\mathfrak{F}}} \mathcal{A}\} \times \bigcup \{Y_b \mid b \in \text{atoms}^{\tilde{\mathfrak{F}}} \mathcal{B}\} = A \times B$ where $A \in \text{up } \mathcal{A}, B \in \text{up } \mathcal{B}$; $K \in \text{up}(\mathcal{A} \times^{\text{RLD}} \mathcal{B})$. □

Theorem 124. $\mathcal{A}_0 \times^{\text{RLD}} \mathcal{B}_0 \cap^{\text{RLD}} \mathcal{A}_1 \times^{\text{RLD}} \mathcal{B}_1 = (\mathcal{A}_0 \cap^{\text{RLD}} \mathcal{A}_1) \times^{\text{RLD}} (\mathcal{B}_0 \cap^{\text{RLD}} \mathcal{B}_1)$ for any $\mathcal{A}_0, \mathcal{A}_1, \mathcal{B}_0, \mathcal{B}_1 \in \tilde{\mathfrak{F}}$.

Proof.

$$\begin{aligned} \mathcal{A}_0 \times^{\text{RLD}} \mathcal{B}_0 \cap^{\text{RLD}} \mathcal{A}_1 \times^{\text{RLD}} \mathcal{B}_1 &= \bigcap^{\text{RLD}} \{P \cap Q \mid P \in \text{up}(\mathcal{A}_0 \times^{\text{RLD}} \mathcal{B}_0), Q \in \text{up}(\mathcal{A}_1 \times^{\text{RLD}} \mathcal{B}_1)\} \\ &= \bigcap^{\text{RLD}} \{A_0 \times B_0 \cap A_1 \times B_1 \mid A_0 \in \text{up } \mathcal{A}_0, B_0 \in \text{up } \mathcal{B}_0, A_1 \in \text{up } \mathcal{A}_1, \\ &\quad B_1 \in \text{up } \mathcal{B}_1\} \\ &= \bigcap^{\text{RLD}} \{(A_0 \cap A_1) \times (B_0 \cap B_1) \mid A_0 \in \text{up } \mathcal{A}_0, B_0 \in \text{up } \mathcal{B}_0, A_1 \in \text{up } \mathcal{A}_1, \\ &\quad B_1 \in \text{up } \mathcal{B}_1\} \\ &= \bigcap^{\text{RLD}} \{K \times L \mid K \in \text{up}(\mathcal{A}_0 \cap \mathcal{A}_1), L \in \text{up}(\mathcal{B}_0 \cap \mathcal{B}_1)\} \\ &= (\mathcal{A}_0 \cap^{\text{RLD}} \mathcal{A}_1) \times^{\text{RLD}} (\mathcal{B}_0 \cap^{\text{RLD}} \mathcal{B}_1). \end{aligned}$$

□

Theorem 125. If $S \in \mathcal{P}\tilde{\mathfrak{F}}^2$ then

$$\bigcap^{\text{RLD}} \{\mathcal{A} \times^{\text{RLD}} \mathcal{B} \mid (\mathcal{A}; \mathcal{B}) \in S\} = \bigcap^{\tilde{\mathfrak{F}}} \text{dom } S \times^{\text{RLD}} \bigcap^{\tilde{\mathfrak{F}}} \text{im } S.$$

Proof. Let $\mathcal{P} = \bigcap^{\tilde{\mathfrak{F}}} \text{dom } S, \mathcal{Q} = \bigcap^{\tilde{\mathfrak{F}}} \text{im } S; l = \bigcap^{\text{RLD}} \{\mathcal{A} \times^{\text{RLD}} \mathcal{B} \mid (\mathcal{A}; \mathcal{B}) \in S\}$.

$\mathcal{P} \times^{\text{RLD}} \mathcal{Q} \subseteq l$ is obvious.

Let $F \in \text{up}(\mathcal{P} \times^{\text{RLD}} \mathcal{Q})$. Then exist $P \in \text{up } \mathcal{P}$ and $Q \in \text{up } \mathcal{Q}$ such that $F \supseteq P \times Q$.

$P = P_1 \cap \dots \cap P_n$ where $P_i \in \langle \text{up} \rangle \text{dom } S$ and $Q = Q_1 \cap \dots \cap Q_m$ where $Q_i \in \langle \text{up} \rangle \text{im } S$.

$P \times Q = \bigcap_{i,j} (P_i \times Q_j)$.

$P_i \times Q_j \supseteq \mathcal{A} \times^{\text{RLD}} \mathcal{B}$ for some $(\mathcal{A}; \mathcal{B}) \in S$. $P \times Q = \bigcap_{i,j} (P_i \times Q_j) \supseteq l$. $F \in \text{up } l$. □

Conjecture 126. If $\mathcal{A} \in \mathfrak{F}$ then $\mathcal{A} \times^{\text{RLD}}$ is a complete homomorphism of the lattice \mathfrak{F} to a complete sublattice of the lattice RLD , if also $\mathcal{A} \neq \emptyset$ then it is an isomorphism.

Definition 127. I will call a reloid *convex* iff it is a union of direct products.

I will call two filter objects *isomorphic* when the corresponding filters are isomorphic (in the sense defined in [5]).

Theorem 128. The reloid $\{a\} \times^{\text{RLD}} \mathcal{F}$ is isomorphic to the filter object \mathcal{F} for every $a \in \mathcal{U}$.

Proof. Consider $B = \{a\} \times \mathcal{U}$ and $f = \{(x; (a; x)) \mid x \in \mathcal{U}\}$. Then f is a bijection from \mathcal{U} to B .

If $X \in \text{up } \mathcal{F}$ then $\langle f \rangle X \in B$ and $\langle f \rangle X = \{a\} \times X \in \text{up}(\{a\} \times^{\text{RLD}} \mathcal{F})$.

For every $Y \in \text{up}(\{a\} \times^{\text{RLD}} \mathcal{F}) \cap \mathcal{P}B$ we have $Y = \{a\} \times X$ for some $X \in \text{up } \mathcal{F}$ and thus $Y = \langle f \rangle X$.

So $\langle f \rangle|_{\text{up } \mathcal{F} \cap \mathcal{P}B} = \langle f \rangle|_{\text{up } \mathcal{F}}$ is a bijection from $\text{up } \mathcal{F} \cap \mathcal{P}B$ to $\text{up}(\{a\} \times^{\text{RLD}} \mathcal{F}) \cap \mathcal{P}B$.

We have $\text{up } \mathcal{F} \cap \mathcal{P}B$ and $\text{up}(\{a\} \times^{\text{RLD}} \mathcal{F}) \cap \mathcal{P}B$ directly isomorphic and thus $\text{up } \mathcal{F}$ is isomorphic to $\text{up}(\{a\} \times^{\text{RLD}} \mathcal{F})$. \square

4.3 Restricting reloid to a filter object. Domain and image

Definition 129. I call restricting a reloid f to a filter object \mathcal{A} as $f|_{\mathcal{A}} = f \cap^{\text{RLD}} \mathcal{A} \times^{\text{RLD}} \mathcal{U}$.

Definition 130. *Domain* and *image* of a reloid f are defined as follows:

$$\text{dom } f = \bigcap^{\mathfrak{F}} \langle \text{dom} \rangle \text{up } f; \quad \text{im } f = \bigcap^{\mathfrak{F}} \langle \text{im} \rangle \text{up } f.$$

Proposition 131. $f \subseteq \mathcal{A} \times^{\text{RLD}} \mathcal{B} \Leftrightarrow \text{dom } f \subseteq \mathcal{A} \wedge \text{im } f \subseteq \mathcal{B}$.

Proof.

Direct implication. Follows from $\text{dom}(\mathcal{A} \times^{\text{RLD}} \mathcal{B}) \subseteq \mathcal{A} \wedge \text{im}(\mathcal{A} \times^{\text{RLD}} \mathcal{B}) \subseteq \mathcal{B}$.

Reverse implication. $\text{dom } f \subseteq \mathcal{A} \Leftrightarrow \forall A \in \text{up } \mathcal{A} \exists F \in \text{up } f: \text{dom } F \subseteq A$. Analogously

$$\text{im } f \subseteq \mathcal{B} \Leftrightarrow \forall B \in \text{up } \mathcal{B} \exists G \in \text{up } f: \text{im } G \subseteq B.$$

Let $\text{dom } f \subseteq \mathcal{A} \wedge \text{im } f \subseteq \mathcal{B}$, $A \in \text{up } \mathcal{A}$, $B \in \text{up } \mathcal{B}$. Then exist $F \in \text{up } f$, $G \in \text{up } f$ such that $\text{dom } F \subseteq A \wedge \text{im } G \subseteq B$. Consequently $F \cap G \in \text{up } f$, $\text{dom}(F \cap G) \subseteq A$, $\text{im}(F \cap G) \subseteq B$ that is $F \cap G \subseteq A \times B$. We have exists $H \in \text{up } f$ such that $H \subseteq A \times B$ for any $A \in \text{up } \mathcal{A}$, $B \in \text{up } \mathcal{B}$. So $f \subseteq \mathcal{A} \times^{\text{RLD}} \mathcal{B}$. \square

Definition 132. I call *identity reloid* for a filter object \mathcal{A} the reloid $I_{\mathcal{A}} \stackrel{\text{def}}{=} (=)|_{\mathcal{A}}$.

Theorem 133. $I_{\mathcal{A}} = \bigcap^{\mathfrak{F}} \{I_A \mid A \in \text{up } \mathcal{A}\}$ where I_A is the identity relation on a set A .

Proof. Let $K \in \text{up } \bigcap^{\mathfrak{F}} \{I_A \mid A \in \text{up } \mathcal{A}\}$, then exists $A \in \text{up } \mathcal{A}$ such that $K \supseteq I_A$. Then $(=)|_{\mathcal{A}} = (=) \cap^{\text{RLD}} \mathcal{A} \times \mathcal{U} \subseteq (=) \cap A \times \mathcal{U} = I|_A \subseteq K$; $K \in \text{up } I_{\mathcal{A}}$. Reversely let $K \in \text{up } I_{\mathcal{A}} = \text{up}((=) \cap^{\text{RLD}} \mathcal{A} \times^{\text{RLD}} \mathcal{U})$, then exists $A \in \text{up } \mathcal{A}$ such that $K \in \text{up}((=) \cap A \times \mathcal{U}) = \text{up } I_A \subseteq \text{up } I_{\mathcal{A}}$. \square

Proposition 134. $I_{\mathcal{A}}^{-1} = I_{\mathcal{A}}$.

Proof. Follows from the previous theorem. \square

Theorem 135. $f|_{\mathcal{A}} = f \circ I_{\mathcal{A}}$ for any reloid f and filter object \mathcal{A} .

Proof. We need to prove that $f \cap^{\text{RLD}} \mathcal{A} \times \mathcal{U} = f \circ \bigcap^{\text{RLD}} \{I_A \mid A \in \text{up } \mathcal{A}\}$. $f \circ \bigcap^{\text{RLD}} \{I_A \mid A \in \text{up } \mathcal{A}\} = \bigcap^{\text{RLD}} \{F \circ I_A \mid F \in \text{up } f, A \in \text{up } \mathcal{A}\} = \bigcap^{\text{RLD}} \{F|_A \mid F \in \text{up } f, A \in \text{up } \mathcal{A}\} = \bigcap^{\text{RLD}} \{F \cap A \times \mathcal{U} \mid F \in \text{up } f, A \in \text{up } \mathcal{A}\} = \bigcap^{\text{RLD}} \{F \mid F \in \text{up } f\} \cap \bigcap^{\text{RLD}} \{A \times \mathcal{U} \mid A \in \text{up } \mathcal{A}\} = f \cap^{\text{RLD}} \mathcal{A} \times^{\text{RLD}} \mathcal{U}$. \square

Theorem 136. $(g \circ f)|_{\mathcal{A}} = g \circ (f|_{\mathcal{A}})$ for any reloids f and g and filter object \mathcal{A} .

Proof. $(g \circ f)|_{\mathcal{A}} = (g \circ f) \circ I_{\mathcal{A}} = g \circ (f \circ I_{\mathcal{A}}) = g \circ (f|_{\mathcal{A}})$. \square

Theorem 137. $f \cap^{\text{RLD}} \mathcal{A} \times^{\text{RLD}} \mathcal{B} = I_{\mathcal{B}} \circ f \circ I_{\mathcal{A}}$ for any reloid f and filter objects \mathcal{A} and \mathcal{B} .

Proof. $f \cap^{\text{RLD}} \mathcal{A} \times^{\text{RLD}} \mathcal{B} = f \cap^{\text{RLD}} \mathcal{A} \times^{\text{RLD}} \mathcal{U} \cap^{\text{RLD}} \mathcal{U} \times^{\text{RLD}} \mathcal{B} = f|_{\mathcal{A}} \cap^{\text{RLD}} \mathcal{U} \times \mathcal{B} = f \circ I_{\mathcal{A}} \cap^{\text{RLD}} \mathcal{U} \times \mathcal{B} = ((f \circ I_{\mathcal{A}})^{-1} \cap^{\text{RLD}} (\mathcal{U} \times^{\text{RLD}} \mathcal{B})^{-1})^{-1} = (I_{\mathcal{A}} \circ f^{-1} \cap^{\text{RLD}} \mathcal{B} \times^{\text{RLD}} \mathcal{U})^{-1} = (I_{\mathcal{A}} \circ f^{-1} \circ I_{\mathcal{B}})^{-1} = I_{\mathcal{B}} \circ f \circ I_{\mathcal{A}}$. \square

4.4 Category of reloids

I will define the category RLD of reloids:

- The set of objects is \mathfrak{F} .
- The set of morphisms from a filter object \mathcal{A} to a filter object \mathcal{B} is the set of triples $(f; \mathcal{A}; \mathcal{B})$ where f is a reloid such that $\text{dom } f \subseteq \mathcal{A}$, $\text{im } f \subseteq \mathcal{B}$.
- Composition of morphisms is defined in the natural way.
- Identity morphism of a filter object \mathcal{A} is $(I_{\mathcal{A}}; \mathcal{A}; \mathcal{A})$.

To prove that it is really a category is trivial.

4.4.1 Monovalued reloids

Following the idea of definition of monovalued morphism let's call *monovalued* such a reloid f that $f \circ f^{-1} \subseteq I_{\text{im } f}$.

Obvious 138. A morphism $(f; \mathcal{A}; \mathcal{B})$ of the category of reloids is monovalued iff the reloid f is monovalued.

Conjecture 139. If a reloid is monovalued then it is a monovalued function restricted to some filter object.

Conjecture 140. A reloid f is monovalued iff $\forall g \in \text{RLD}: (g \subseteq f \Rightarrow \exists \mathcal{A} \in \mathfrak{F}: g = f|_{\mathcal{A}})$.

Conjecture 141. A monovalued reloid restricted to an atomic filter object is atomic or empty.

A weaker conjecture:

Conjecture 142. A (monovalued) function restricted to an atomic filter object is atomic or empty.

5 Relationships of funcoids and reloids

5.1 Funcoid induced by a reloid

Every reloid f induces a funcoid $(\text{FCD})f$ by the following formulas:

$$\begin{aligned} \mathcal{X}[(\text{FCD})f]\mathcal{Y} &\Leftrightarrow \forall F \in \text{up } f: \mathcal{X}[F]\mathcal{Y} \\ \langle (\text{FCD})f \rangle \mathcal{X} &= \bigcap^{\mathfrak{F}} \{ \langle F \rangle \mathcal{X} \mid F \in \text{up } f \}. \end{aligned}$$

We should prove that $(\text{FCD})f$ is really a funcoid. For this purpose we will additionally define

$$\langle (\text{FCD})f^{-1} \rangle \mathcal{Y} = \bigcap^{\mathfrak{F}} \{ \langle F^{-1} \rangle \mathcal{Y} \mid F \in \text{up } f \}.$$

Proof. We need to prove that

$$\mathcal{X}[(\text{FCD})f]\mathcal{Y} \Leftrightarrow \mathcal{Y} \cap^{\mathfrak{F}} \langle (\text{FCD})f \rangle \mathcal{X} \neq \emptyset \Leftrightarrow \mathcal{X} \cap^{\mathfrak{F}} \langle (\text{FCD})f^{-1} \rangle \mathcal{Y} \neq \emptyset.$$

The above formula is equivalent to:

$$\forall F \in \text{up } f: \mathcal{X}[F]\mathcal{Y} \Leftrightarrow \mathcal{Y} \cap^{\mathfrak{F}} \bigcap^{\mathfrak{F}} \{ \langle F \rangle \mathcal{X} \mid F \in \text{up } f \} \neq \emptyset \Leftrightarrow \mathcal{X} \cap^{\mathfrak{F}} \bigcap^{\mathfrak{F}} \{ \langle F^{-1} \rangle \mathcal{Y} \mid F \in \text{up } f \} \neq \emptyset.$$

We have $\mathcal{Y} \cap^{\mathfrak{F}} \bigcap^{\mathfrak{F}} \{ \langle F \rangle \mathcal{X} \mid F \in \text{up } f \} = \bigcap^{\mathfrak{F}} \{ \mathcal{Y} \cap^{\mathfrak{F}} \langle F \rangle \mathcal{X} \mid F \in \text{up } f \}$.

Let's denote $W = \{ \mathcal{Y} \cap^{\mathfrak{F}} \langle F \rangle \mathcal{X} \mid F \in \text{up } f \}$.

We need to prove that $\bigcap^{\mathfrak{F}} W \neq \emptyset \Leftrightarrow \forall F \in \text{up } f: \mathcal{Y} \cap^{\mathfrak{F}} \langle F \rangle \mathcal{X} \neq \emptyset$. (The rest follows from symmetry.)

Let's prove that W is a generalized filter base. [TODO: Simplify this proof.] Let $\mathcal{A}, \mathcal{B} \in W$ that is $\mathcal{A} = \mathcal{Y} \cap^{\mathfrak{F}} \langle P \rangle \mathcal{X}$, $\mathcal{B} = \mathcal{Y} \cap^{\mathfrak{F}} \langle Q \rangle \mathcal{X}$ where $P, Q \in \text{up } f$. Then for $\mathcal{C} = \mathcal{Y} \cap^{\mathfrak{F}} \langle P \cap Q \rangle \mathcal{X}$ is true both $\mathcal{C} \in W$ and $\mathcal{C} \subseteq \mathcal{A}, \mathcal{B}$. So W is a generalized filter base.

From this by the corollary 4 follows that $\bigcap^{\mathfrak{F}} W \neq \emptyset \Leftrightarrow \emptyset \notin W \Leftrightarrow \forall F \in \text{up } f: \mathcal{Y} \cap^{\mathfrak{F}} \langle F \rangle \mathcal{X} \neq \emptyset$. \square

Theorem 143. $\mathcal{X}[(\text{FCD})f]\mathcal{Y} \Leftrightarrow \mathcal{X} \times^{\text{RLD}} \mathcal{Y} \cap f \neq \emptyset$ for any $\mathcal{X}, \mathcal{Y} \in \mathfrak{F}$ and $f \in \text{RLD}$.

Proof.

$$\begin{aligned} \mathcal{X} \times^{\text{RLD}} \mathcal{Y} \cap^{\text{RLD}} f \neq \emptyset &\Leftrightarrow \forall P \in \text{up}(\mathcal{X} \times^{\text{RLD}} \mathcal{Y}): P \cap^{\text{RLD}} f \neq \emptyset \\ &\Leftrightarrow \forall F \in \text{up } f, P \in \text{up}(\mathcal{X} \times^{\text{RLD}} \mathcal{Y}): P \cap F \neq \emptyset \\ &\Leftrightarrow \forall F \in \text{up } f, X \in \text{up } \mathcal{X}, Y \in \text{up } \mathcal{Y}: X \times^{\text{RLD}} Y \cap^{\text{RLD}} F \neq \emptyset \\ &\Leftrightarrow \forall F \in \text{up } f, X \in \text{up } \mathcal{X}, Y \in \text{up } \mathcal{Y}: X[F]Y \\ &\Leftrightarrow \forall F \in \text{up } f: \mathcal{X}[F]\mathcal{Y} \\ &\Leftrightarrow \mathcal{X}[(\text{FCD})f]\mathcal{Y}. \end{aligned}$$

\square

Theorem 144. $(\text{FCD})f = \bigcap^{\text{FCD}} \text{up } f$ for any reloid f .

Proof. Let a is an atomic filter object.

$((\text{FCD})f)a = \bigcap^{\mathfrak{F}} \{ \langle F \rangle a \mid F \in \text{up } f \}$ by the definition of (FCD) .

$\langle \bigcap^{\text{FCD}} \text{up } f \rangle a = \bigcap^{\mathfrak{F}} \{ \langle F \rangle a \mid F \in \text{up } f \}$ by the theorem 53.

So $\langle (\text{FCD})f \rangle a = \langle \bigcap^{\text{FCD}} \text{up } f \rangle a$ for any atomic filter object a . \square

Lemma 145. $\langle g \rangle \bigcap^{\mathfrak{F}} S = \bigcap^{\mathfrak{F}} \langle \langle g \rangle \rangle S$ if g is a funcoid and S is a filter base.

Proof. $\text{up} \bigcap^{\mathfrak{F}} S = \bigcup \langle \text{up} \rangle S$ by the theorem 3.

$\langle g \rangle \bigcap^{\mathfrak{F}} S = \bigcap^{\mathfrak{F}} \langle \langle g \rangle \rangle \text{up} \bigcap^{\mathfrak{F}} S$ by the theorem 33.

$\bigcap^{\mathfrak{F}} \langle \langle g \rangle \rangle \text{up} \bigcap^{\mathfrak{F}} S = \bigcap^{\mathfrak{F}} \langle \langle g \rangle \rangle \bigcup \langle \text{up} \rangle S$.

Easy to see that $\bigcap^{\mathfrak{F}} \langle \langle g \rangle \rangle \bigcup \langle \text{up} \rangle S = \bigcap^{\mathfrak{F}} \langle \langle g \rangle \rangle S$ because $S \subseteq \bigcup \langle \text{up} \rangle S$.

Combining these equalities we produce $\langle g \rangle \bigcap^{\mathfrak{F}} S = \bigcap^{\mathfrak{F}} \langle \langle g \rangle \rangle S$. \square

Lemma 146. For two sets of binary relations S and T and a set A

$$\bigcap^{\mathfrak{F}} S = \bigcap^{\mathfrak{F}} T \Rightarrow \bigcap^{\mathfrak{F}} \{ \langle F \rangle A \mid F \in S \} = \bigcap^{\mathfrak{F}} \{ \langle G \rangle A \mid G \in T \}$$

Proof. Let $\bigcap^{\mathfrak{F}} S = \bigcap^{\mathfrak{F}} T$. Suppose $X \in \bigcap^{\mathfrak{F}} \{ \langle F \rangle A \mid F \in S \}$. Then $X' \in \{ \langle F \rangle A \mid F \in S \}$ where $X \supseteq X'$. That is $X' = \langle F \rangle A$ for some $F \in S$. There exists $G \in T$ such that $G \subseteq F$. So $Y' = \langle G \rangle A \subseteq X' \subseteq X$. $Y' \in \{ \langle G \rangle A \mid G \in T \}$; $Y' \in \text{up} \bigcap^{\mathfrak{F}} \{ \langle G \rangle A \mid G \in T \}$; $X \in \text{up} \bigcap^{\mathfrak{F}} \{ \langle G \rangle A \mid G \in T \}$. The reverse is symmetric. \square

Theorem 147. $(\text{FCD})(g \circ f) = ((\text{FCD})g) \circ ((\text{FCD})f)$ for any reloids f and g .

Proof.

$$\begin{aligned} \langle (\text{FCD})(g \circ f) \rangle X &= \bigcap^{\mathfrak{F}} \{ \langle H \rangle X \mid H \in \text{up}(g \circ f) \} \\ &= \bigcap^{\mathfrak{F}} \left\{ \langle H \rangle X \mid H \in \text{up} \bigcap^{\text{RLD}} \{ G \circ F \mid F \in \text{up } f, G \in \text{up } g \} \right\}. \end{aligned}$$

Obviously

$$\bigcap^{\text{RLD}}\{G \circ F \mid F \in \text{up } f, G \in \text{up } g\} = \bigcap^{\text{RLD}} \text{up} \bigcap^{\text{RLD}}\{G \circ F \mid F \in \text{up } f, G \in \text{up } g\};$$

from this by the lemma 146

$$\bigcap^{\mathfrak{F}}\{\langle H \rangle X \mid H \in \text{up} \bigcap^{\text{RLD}}\{G \circ F \mid F \in \text{up } f, G \in \text{up } g\}\} = \bigcap^{\mathfrak{F}}\{\langle G \circ F \rangle X \mid F \in \text{up } f, G \in \text{up } g\}.$$

On the other side

$$\begin{aligned} \langle ((\text{FCD})g) \circ ((\text{FCD})f) \rangle X &= \langle (\text{FCD})g \rangle \langle (\text{FCD})f \rangle X \\ &= \langle (\text{FCD})g \rangle \bigcap^{\mathfrak{F}}\{\langle F \rangle X \mid F \in \text{up } f\} \\ &= \bigcap^{\mathfrak{F}}\{\langle G \rangle \bigcap^{\mathfrak{F}}\{\langle F \rangle X \mid F \in \text{up } f\} \mid G \in \text{up } g\}. \end{aligned}$$

Let's prove that $\{\langle F \rangle X \mid F \in \text{up } f\}$ is a filter base. If $A, B \in \{\langle F \rangle X \mid F \in \text{up } f\}$ then $A = \langle F_1 \rangle X$ and $B = \langle F_2 \rangle X$ where $F_1, F_2 \in \text{up } f$. $A \cap B \supseteq \langle F_1 \cap F_2 \rangle X \in \{\langle F \rangle X \mid F \in \text{up } f\}$. So $\{\langle F \rangle X \mid F \in \text{up } f\}$ is really a filter base.

By the lemma 145 $\langle G \rangle \bigcap^{\mathfrak{F}}\{\langle F \rangle X \mid F \in \text{up } f\} = \bigcap^{\mathfrak{F}}\{\langle G \rangle \langle F \rangle X \mid F \in \text{up } f\}$. So continuing the above equalities,

$$\begin{aligned} \langle ((\text{FCD})g) \circ ((\text{FCD})f) \rangle X &= \bigcap^{\mathfrak{F}}\{\bigcap^{\mathfrak{F}}\{\langle G \rangle \langle F \rangle X \mid F \in \text{up } f\} \mid G \in \text{up } g\} \\ &= \bigcap^{\mathfrak{F}}\{\langle G \rangle \langle F \rangle X \mid F \in \text{up } f, G \in \text{up } g\} \\ &= \bigcap^{\mathfrak{F}}\{\langle G \circ F \rangle X \mid F \in \text{up } f, G \in \text{up } g\}. \end{aligned}$$

Combining these equalities we get $\langle (\text{FCD})(g \circ f) \rangle X = \langle ((\text{FCD})g) \circ ((\text{FCD})f) \rangle X$ for any set X . \square

5.2 Reloids induced by funcoid

Every funcoid f induces a reloid in two ways, intersection of *outward* relations and union of *inward* direct products of filter objects:

$$\begin{aligned} (\text{RLD})_{\text{out}} f &\stackrel{\text{def}}{=} \bigcap^{\text{RLD}}\{F \mid F \in \mathcal{P}\mathcal{U}^2, F \supseteq f\} \\ (\text{RLD})_{\text{in}} f &\stackrel{\text{def}}{=} \bigcup^{\text{RLD}}\{\mathcal{A} \times^{\text{RLD}} \mathcal{B} \mid \mathcal{A}, \mathcal{B} \in \mathfrak{F}, \mathcal{A} \times^{\text{FCD}} \mathcal{B} \subseteq f\} \end{aligned}$$

Proposition 148. $\text{up}(\text{RLD})_{\text{out}} f = \{F \mid F \in \mathcal{P}\mathcal{U}^2, F \supseteq f\}$.

Proof. Easy to prove. \square

Theorem 149. $(\text{RLD})_{\text{in}} f = \bigcup^{\text{RLD}}\{a \times^{\text{RLD}} b \mid a, b \in \text{atoms}^{\mathfrak{F}}\mathcal{U}, a \times^{\text{FCD}} b \subseteq f\}$.

Proof. Follows from the theorem 123. \square

Lemma 150. $F \in \text{up}(\text{RLD})_{\text{in}} f \Leftrightarrow \forall a, b \in \text{atoms}^{\mathfrak{F}}\mathcal{U}: (a[f]b \Rightarrow F \supseteq a \times^{\text{RLD}} b)$ for a funcoid f .

Proof.

$$\begin{aligned} F \in \text{up}(\text{RLD})_{\text{in}} f &\Leftrightarrow F \in \text{up} \bigcup^{\mathfrak{F}}\{a \times^{\text{RLD}} b \mid a, b \in \text{atoms}^{\mathfrak{F}}\mathcal{U}, a \times^{\text{FCD}} b \subseteq f\} \\ &\Leftrightarrow \forall a, b \in \text{atoms}^{\mathfrak{F}}\mathcal{U}: (a \times^{\text{FCD}} b \subseteq f \Rightarrow F \in \text{up}(a \times^{\text{RLD}} b)) \\ &\Leftrightarrow \forall a, b \in \text{atoms}^{\mathfrak{F}}\mathcal{U}: (a \times^{\text{FCD}} b \cap^{\text{FCD}} f \neq \emptyset \Rightarrow F \supseteq a \times^{\text{RLD}} b) \\ &\Leftrightarrow \forall a, b \in \text{atoms}^{\mathfrak{F}}\mathcal{U}: (a[f]b \Rightarrow F \supseteq a \times^{\text{RLD}} b). \end{aligned}$$

\square

Surprisingly a funcoid is greater inward than outward:

Theorem 151. $(\text{RLD})_{\text{out}} f \subseteq (\text{RLD})_{\text{in}} f$ for a funcoid f .

Proof. We need to prove

$$\bigcap^{\text{RLD}} \{F \mid F \in \mathcal{P}\mathcal{U}^2, F \supseteq f\} \subseteq \bigcup^{\text{RLD}} \{\mathcal{A} \times^{\text{RLD}} \mathcal{B} \mid \mathcal{A}, \mathcal{B} \in \mathfrak{F}, \mathcal{A} \times^{\text{FCD}} \mathcal{B} \subseteq f\}.$$

Let

$$K \in \text{up} \bigcup^{\mathfrak{F}} \{\mathcal{A} \times^{\text{RLD}} \mathcal{B} \mid \mathcal{A}, \mathcal{B} \in \mathfrak{F}, \mathcal{A} \times^{\text{FCD}} \mathcal{B} \subseteq f\}.$$

Then

$$\begin{aligned} K &= \bigcup \{X_{\mathcal{A}} \times Y_{\mathcal{B}} \mid \mathcal{A}, \mathcal{B} \in \mathfrak{F}, \mathcal{A} \times^{\text{FCD}} \mathcal{B} \subseteq f\} \\ &= \bigcup^{\text{RLD}} \{X_{\mathcal{A}} \times Y_{\mathcal{B}} \mid \mathcal{A}, \mathcal{B} \in \mathfrak{F}, \mathcal{A} \times^{\text{FCD}} \mathcal{B} \subseteq f\} \\ &\supseteq f \end{aligned}$$

where $X_{\mathcal{A}} \in \text{up} \mathcal{A}$, $Y_{\mathcal{B}} \in \text{up} \mathcal{B}$. $K \in \text{up} \bigcap^{\text{RLD}} \{F \mid F \in \mathcal{P}\mathcal{U}^2, F \supseteq f\}$. □

Theorem 152. $(\text{FCD})(\text{RLD})_{\text{out}} f = f$ for any functor f .

Proof. $\text{up}(\text{RLD})_{\text{out}} f = \{F \mid F \in \mathcal{P}\mathcal{U}^2, F \supseteq f\}$

$$(\text{FCD})(\text{RLD})_{\text{out}} f = \bigcap^{\text{FCD}} \text{up}(\text{RLD})_{\text{out}} f = \bigcap^{\text{FCD}} \{F \mid F \in \mathcal{P}\mathcal{U}^2, F \supseteq f\}.$$

$$\bigcap^{\text{FCD}} \{F \mid F \in \mathcal{P}\mathcal{U}^2, F \supseteq f\} = f \text{ by the theorem 87. So } (\text{FCD})(\text{RLD})_{\text{out}} f = f. \quad \square$$

Conjecture 153. $(\text{FCD})(\text{RLD})_{\text{in}} f = f$ for any functor f .

Conjecture 154. For any functor f and reloid g

$$(\text{RLD})_{\text{out}} f \subseteq g \subseteq (\text{RLD})_{\text{in}} f \Leftrightarrow (\text{FCD})g = f.$$

Conjecture 155. For a convex reloid f

1. $(\text{RLD})_{\text{out}}(\text{FCD})f = f$;
2. $(\text{RLD})_{\text{in}}(\text{FCD})f = f$.

6 Continuous morphisms

This section will use the apparatus from the section “Partially ordered dagger categories”.

6.1 Traditional definitions of continuity

6.1.1 Pre-topology

Let μ and ν are functors representing some pre-topologies. By definition a function f is continuous map from μ to ν in point a iff

$$\forall \epsilon \in \text{up} \langle \nu \rangle f a \exists \delta \in \text{up} \langle \mu \rangle \{a\}: \langle f \rangle \delta \subseteq \epsilon.$$

Equivalently transforming this formula we get:

$$\begin{aligned} \forall \epsilon \in \text{up} \langle \nu \rangle f a: \langle f \rangle \langle \mu \rangle \{a\} &\subseteq \epsilon; \\ \langle f \rangle \langle \mu \rangle \{a\} &\subseteq \langle \nu \rangle f a; \\ \langle f \rangle \langle \mu \rangle \{a\} &\subseteq \langle \nu \rangle \langle f \rangle \{a\}. \end{aligned}$$

f is a continuous map from μ to ν in every point of its domain iff $\langle f \rangle \langle \mu \rangle \subseteq \langle \nu \rangle \langle f \rangle$ what is equivalent to $f \circ \mu \subseteq \nu \circ f$.

6.1.2 Proximity spaces

Let μ and ν are proximity (nearness) spaces (which I consider a special case of functors). By definition a function f is a nearness-continuous map from μ to ν iff

$$\forall X, Y \in \mathcal{P}\mathcal{U}: (X[\mu]Y \Rightarrow (\langle f \rangle X)[\nu](\langle f \rangle Y)).$$

Equivalently transforming this formula we get:

$$\begin{aligned}
& \forall X, Y \in \mathcal{P}\mathcal{U}: (X[\mu]Y \Rightarrow \langle f \rangle Y \cap \langle \nu \rangle \langle f \rangle X \neq \emptyset); \\
& \forall X, Y \in \mathcal{P}\mathcal{U}: (X[\mu]Y \Rightarrow \langle f \rangle Y \cap \langle \nu \circ f \rangle X \neq \emptyset); \\
& \quad \forall X, Y \in \mathcal{P}\mathcal{U}: (X[\mu]Y \Rightarrow X[\nu \circ f] \langle f \rangle Y); \\
& \quad \forall X, Y \in \mathcal{P}\mathcal{U}: (X[\mu]Y \Rightarrow \langle f \rangle Y[(\nu \circ f)^{-1}]X); \\
& \quad \forall X, Y \in \mathcal{P}\mathcal{U}: (X[\mu]Y \Rightarrow \langle f \rangle Y[f^{-1} \circ \nu^{-1}]X); \\
& \forall X, Y \in \mathcal{P}\mathcal{U}: (X[\mu]Y \Rightarrow X \cap \langle f^{-1} \circ \nu^{-1} \rangle \langle f \rangle Y \neq \emptyset); \\
& \forall X, Y \in \mathcal{P}\mathcal{U}: (X[\mu]Y \Rightarrow X \cap \langle f^{-1} \circ \nu^{-1} \circ f \rangle Y \neq \emptyset); \\
& \quad \forall X, Y \in \mathcal{P}\mathcal{U}: (X[\mu]Y \Rightarrow Y[f^{-1} \circ \nu^{-1} \circ f]X); \\
& \quad \forall X, Y \in \mathcal{P}\mathcal{U}: (X[\mu]Y \Rightarrow X[f^{-1} \circ \nu \circ f]Y); \\
& \quad \mu \subseteq f^{-1} \circ \nu \circ f.
\end{aligned}$$

So a function f is nearness-continuous iff $\mu \subseteq f^{-1} \circ \nu \circ f$.

6.1.3 Uniform spaces

Uniform spaces are a special case of reloids.

Let μ and ν are uniform spaces. By definition a function f is a uniformly continuous map from μ to ν iff

$$\forall \epsilon \in \text{up } \nu \exists \delta \in \text{up } \mu \forall (x; y) \in \delta: (fx; fy) \in \epsilon.$$

Equivalently transforming this formula we get:

$$\begin{aligned}
& \forall \epsilon \in \text{up } \nu \exists \delta \in \text{up } \mu \forall (x; y) \in \delta: \{(fx; fy)\} \subseteq \epsilon \\
& \forall \epsilon \in \text{up } \nu \exists \delta \in \text{up } \mu \forall (x; y) \in \delta: f \circ \{(x; y)\} \circ f^{-1} \subseteq \epsilon \\
& \quad \forall \epsilon \in \text{up } \nu \exists \delta \in \text{up } \mu: f \circ \delta \circ f^{-1} \subseteq \epsilon \\
& \quad \forall \epsilon \in \text{up } \nu: f \circ \mu \circ f^{-1} \subseteq \epsilon \\
& \quad f \circ \mu \circ f^{-1} \subseteq \nu.
\end{aligned}$$

So a function f is uniformly continuous iff $f \circ \mu \circ f^{-1} \subseteq \nu$.

6.2 Our three definitions of continuity

I have expressed different kinds of continuity with simple algebraic formulas hiding the complexity of traditional epsilon-delta notation behind a smart algebra. Let's summarize these three algebraic formulas:

Let μ and ν are endomorphisms of some partially ordered precategory. Continuous functions can be defined as these morphisms f of this precategory which conform to the following formula:

$$f \in C(\mu; \nu) \Leftrightarrow f \in \text{Mor}(\text{Ob } \mu; \text{Ob } \nu) \wedge f \circ \mu \subseteq \nu \circ f.$$

If the precategory is a partially ordered dagger precategory then continuity also can be defined in two other ways:

$$\begin{aligned}
f \in C'(\mu; \nu) & \Leftrightarrow f \in \text{Mor}(\text{Ob } \mu; \text{Ob } \nu) \wedge \mu \subseteq f^\dagger \circ \nu \circ f; \\
f \in C''(\mu; \nu) & \Leftrightarrow f \in \text{Mor}(\text{Ob } \mu; \text{Ob } \nu) \wedge f \circ \mu \circ f^\dagger \subseteq \nu.
\end{aligned}$$

Remark 156. In the examples about funcoids and reloids the “dagger functor” is the inverse of a funcoid or reloid, that is $f^\dagger = f^{-1}$.

Proposition 157. Every of these three definitions of continuity forms a sub-precategory (sub-category if the original precategory is a category).

Proof.

C. Let $f \in C(\mu; \nu)$, $g \in C(\nu; \pi)$. Then $f \circ \mu \subseteq \nu \circ f$, $g \circ \nu \subseteq \pi \circ g$; $g \circ f \circ \mu \subseteq g \circ \nu \circ f \subseteq \pi \circ g \circ f$. So $g \circ f \in C(\mu; \pi)$. $1_{\text{Ob } \mu} \in C(\mu; \mu)$ is obvious.

C' . Let $f \in C'(\mu; \nu)$, $g \in C'(\nu; \pi)$. Then $\mu \subseteq f^\dagger \circ \nu \circ f$, $\nu \subseteq g^\dagger \circ \pi \circ g$;

$$\mu \subseteq f^\dagger \circ g^\dagger \circ \pi \circ g \circ f; \quad \mu \subseteq (g \circ f)^\dagger \circ \pi \circ (g \circ f).$$

So $g \circ f \in C'(\mu; \pi)$. $1_{\text{Ob } \mu} \in C'(\mu; \mu)$ is obvious.

C'' . Let $f \in C''(\mu; \nu)$, $g \in C''(\nu; \pi)$. Then $f \circ \mu \circ f^\dagger \subseteq \nu$, $g \circ \nu \circ g^\dagger \subseteq \pi$;

$$g \circ f \circ \mu \circ f^\dagger \circ g^\dagger \subseteq \pi; \quad (g \circ f) \circ \mu \circ (g \circ f)^\dagger \subseteq \pi.$$

So $g \circ f \in C''(\mu; \pi)$. $1_{\text{Ob } \mu} \in C''(\mu; \mu)$ is obvious. \square

Proposition 158. For a monovalued morphism f of a partially ordered dagger category and its endomorphisms μ and ν

$$f \in C'(\mu; \nu) \Rightarrow f \in C(\mu; \nu) \Rightarrow f \in C''(\mu; \nu).$$

Proof. Let $f \in C'(\mu; \nu)$. Then $\mu \subseteq f^\dagger \circ \nu \circ f$; $f \circ \mu \subseteq f \circ f^\dagger \circ \nu \circ f \subseteq 1_{\text{Dst } f} \circ \nu \circ f = \nu \circ f$; $f \in C(\mu; \nu)$.

Let $f \in C(\mu; \nu)$. Then $f \circ \mu \subseteq \nu \circ f$; $f \circ \mu \circ f^\dagger \subseteq \nu \circ f \circ f^\dagger \subseteq \nu \circ 1_{\text{Dst } f} = \nu$; $f \in C''(\mu; \nu)$. \square

Proposition 159. For an entirely defined morphism f of a partially ordered dagger category and its endomorphisms μ and ν

$$f \in C''(\mu; \nu) \Rightarrow f \in C(\mu; \nu) \Rightarrow f \in C'(\mu; \nu).$$

Proof. Let $f \in C''(\mu; \nu)$. Then $f \circ \mu \circ f^\dagger \subseteq \nu$; $f \circ \mu \circ f^\dagger \circ f \subseteq \nu \circ f$; $f \circ \mu \circ 1_{\text{Src } f} \subseteq \nu \circ f$; $f \circ \mu \subseteq \nu \circ f$; $f \in C(\mu; \nu)$.

Let $f \in C(\mu; \nu)$. Then $f \circ \mu \subseteq \nu \circ f$; $f^\dagger \circ f \circ \mu \subseteq f^\dagger \circ \nu \circ f$; $1_{\text{Src } f} \circ \mu \subseteq f^\dagger \circ \nu \circ f$; $\mu \subseteq f^\dagger \circ \nu \circ f$; $f \in C'(\mu; \nu)$. \square

For entirely defined monovalued morphisms our three definitions of continuity coincide:

Theorem 160. If f is a monovalued and entirely defined morphism then

$$f \in C'(\mu; \nu) \Leftrightarrow f \in C(\mu; \nu) \Leftrightarrow f \in C''(\mu; \nu).$$

Proof. From two previous propositions. \square

The classical general topology theorem that uniformly continuous function from a uniform space to an other uniform space is near-continuous regarding the proximities generated by the uniformities, generalized for reloids and functors takes the following form:

Theorem 161. If an entirely defined morphism of the category of reloids $f \in C''(\mu; \nu)$ for some endomorphisms μ and ν of the category of reloids, then $(\text{FCD})f \in C'((\text{FCD})\mu; (\text{FCD})\nu)$.

Exercise 1. I leave a simple exercise for the reader to prove the last theorem.

6.3 Continuousness of a restricted morphism

Consider some partially ordered semigroup. (For example it can be the semigroup of functors or semigroup of reloids.) Consider also some lattice (*lattice of objects*). (For example take the lattice of set theoretic filters.)

We will map every object A to *identity element* I_A of the semigroup (for example identity functor or identity reloid). For identity elements we will require

1. $I_A \circ I_B = I_{A \cap B}$;
2. $f \circ I_A \subseteq f$; $I_A \circ f \subseteq f$.

In the case when our semigroup is “dagger” (that is is a dagger precategory) we will require also $(I_A)^\dagger = I_A$.

We can define *restricting* an element f of our semigroup to an object A by the formula $f|_A = f \circ I_A$.

We can define *rectangular restricting* an element μ of our semigroup to objects A and B as $I_B \circ \mu \circ I_A$. Optionally we can define direct product $A \times B$ of two objects by the formula (true for functors and for relicts):

$$\mu \cap (A \times B) = I_B \circ \mu \circ I_A.$$

Square restricting of an element μ to an object A is a special case of rectangular restricting and is defined by the formula $I_A \circ \mu \circ I_A$ (or by the formula $\mu \cap (A \times A)$).

Theorem 162. For any elements f, μ, ν of our semigroup and an object A

1. $f \in C(\mu; \nu) \Rightarrow f|_A \in C(I_A \circ \mu \circ I_A; \nu)$;
2. $f \in C'(\mu; \nu) \Rightarrow f|_A \in C'(I_A \circ \mu \circ I_A; \nu)$;
3. $f \in C''(\mu; \nu) \Rightarrow f|_A \in C''(I_A \circ \mu \circ I_A; \nu)$.

(Two last items are true for the case when our semigroup is with inverses.)

Proof.

1. $f|_A \in C(I_A \circ \mu \circ I_A; \nu) \Leftrightarrow f|_A \circ I_A \circ \mu \circ I_A \subseteq \nu \circ f|_A \Leftrightarrow f \circ I_A \circ I_A \circ \mu \circ I_A \subseteq \nu \circ f|_A \Leftrightarrow f \circ I_A \circ \mu \circ I_A \subseteq \nu \circ f \circ I_A \Leftrightarrow f \circ I_A \circ \mu \subseteq \nu \circ f \Leftrightarrow f \circ \mu \subseteq \nu \circ f \Leftrightarrow f \in C(\mu; \nu)$.
2. $f|_A \in C'(I_A \circ \mu \circ I_A; \nu) \Leftrightarrow I_A \circ \mu \circ I_A \subseteq (f|_A)^\dagger \circ \nu \circ f|_A \Leftrightarrow I_A \circ \mu \circ I_A \subseteq (f \circ I_A)^\dagger \circ \nu \circ f \circ I_A \Leftrightarrow I_A \circ \mu \circ I_A \subseteq I_A \circ f^\dagger \circ \nu \circ f \circ I_A \Leftrightarrow \mu \subseteq f^\dagger \circ \nu \circ f \Leftrightarrow f \in C'(\mu; \nu)$.
3. $f|_A \in C''(I_A \circ \mu \circ I_A; \nu) \Leftrightarrow f|_A \circ I_A \circ \mu \circ I_A \circ (f|_A)^\dagger \subseteq \nu \Leftrightarrow f \circ I_A \circ I_A \circ \mu \circ I_A \circ I_A \circ f^\dagger \subseteq \nu \Leftrightarrow f \circ I_A \circ \mu \circ I_A \circ f^\dagger \subseteq \nu \Leftrightarrow f \circ \mu \circ f^\dagger \subseteq \nu \Leftrightarrow f \in C''(\mu; \nu)$. \square

7 Postface

7.1 Misc

I deem that now two most important research topics in Algebraic General Topology are:

- to solve the open problems mentioned in this work;
- define and research compactness of functors.

Also a future research topic are n -ary (where n is an ordinal, or more generally an index set) functors and relicts (plain functors and relicts are binary by analogy with binary relations).

7.2 Pointfree functors and relicts

I have set wiki site <http://functors.wikidot.com> to write on that site the pointfree variant of the theory of functors and relicts (that is generalized functors on arbitrary lattices rather than functors on a lattice of sets as in this work).

However I consider for me research of pointfree functors and pointfree relicts a low priority project. (There are yet enough research topics in the point-set topology and I don't want to meddle into pointfree topology in foreseeable future.)

The work about pointfree functors and relicts seems being largely technical and boring. Pointfree theory of functors and relicts seems being a trivial generalization of the theory of point-set functors and relicts. It is not similar to the traditional pointfree topology which is not an obvious generalization of point-set topology.

But if someone indeed wishes to treat pointfree functors, please use the above mentioned wiki.

Appendix A Some counter-examples

[TODO: More counter-examples similar to examples in [5].]

Theorem 163. For a f.o. a we have $a \times^{\text{RLD}} a \subseteq (=)|_{\mathcal{U}}$ only in the case if $a = \emptyset$ or a is a trivial atomic f.o. (that is an one-element set).

Proof. If $a \times^{\text{RLD}} a \subseteq (=)|_{\mathcal{U}}$ then exists $m \in \text{up}(a \times^{\text{RLD}} a)$ such that $m \subseteq (=)|_{\mathcal{U}}$. Consequently exist $A, B \in \text{up } a$ such that $A \times B \subseteq (=)|_{\mathcal{U}}$ what is possible only in the case when $A = B = a$ is an one-element set or empty set. \square

Corollary 164. Direct product (in the sense of reلودs) of non-trivial atomic filter objects is non-atomic.

Proof. Obviously $(a \times^{\text{RLD}} a) \cap^{\text{RLD}} (=)|_{\mathcal{U}} \neq \emptyset$ and $(a \times^{\text{RLD}} a) \cap^{\text{RLD}} (=)|_{\mathcal{U}} \subset a \times^{\text{RLD}} a$. \square

Example 165. There exist two atomic reلودs whose composition is non-atomic and non-empty.

Proof. Let a is a non-trivial atomic filter object and $x \in \mathcal{U}$. Then

$$(a \times \{x\}) \circ (\{x\} \times a) = \bigcap^{\mathfrak{F}} \{(A \times \{x\}) \circ (\{x\} \times A) \mid A \in \text{up } a\} = \bigcap^{\mathfrak{F}} \{A \times A \mid A \in \text{up } a\} = a \times a$$

is non-atomic despite of $a \times \{x\}$ and $\{x\} \times a$ are atomic. \square

Example 166. There exists non-monovalued atomic reلود.

Proof. From the previous example follows that the atomic reلود $\{x\} \times a$ is not monovalued. \square

[TODO: Example of $(\text{RLD})_{\text{in}} \neq (\text{RLD})_{\text{out}}$.]

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